

WatchWitch

Hacking the Apple Watch





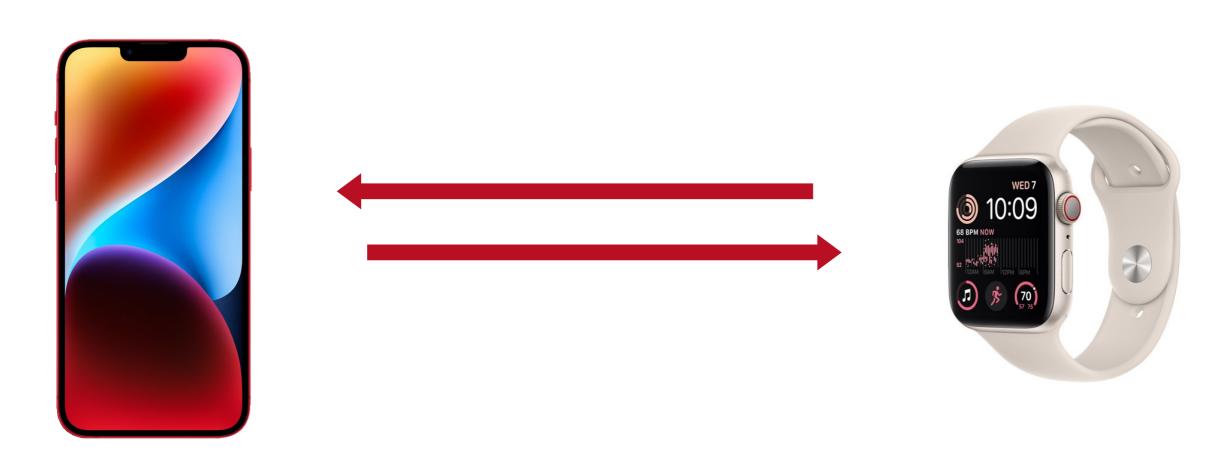
1. Prelude



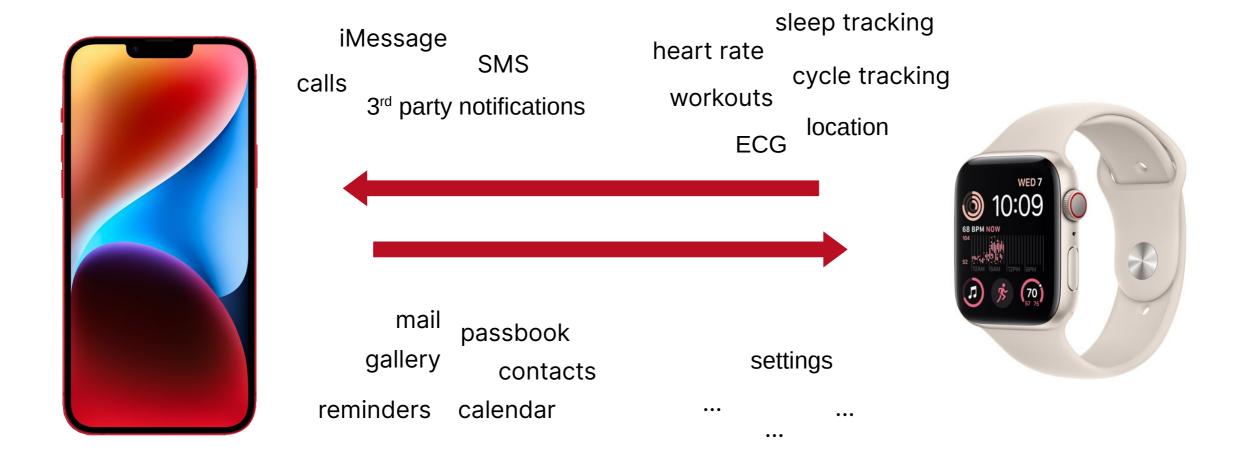














Why should I care?





twitter.com/iamkoby/status/689521611611971588





```
INSERT INTO quantity_samples VALUES(48802,0.42,NULL,NULL);
INSERT INTO quantity_samples VALUES(48803,0.39,NULL,NULL);
INSERT INTO quantity_samples VALUES(48804,0.98,59.0,1);
INSERT INTO quantity_samples VALUES(48824, 0.65, NULL, NULL);
INSERT INTO quantity_samples VALUES(48825,12.2,NULL,NULL);
INSERT INTO quantity_samples VALUES(48858,1.35,81.0,1)
sqlite> SELECT private_classification, average_heart_rate,
hex(voltage_payload) FROM ecg_samples LIMIT 1;
private_classification|average_heart_rate|voltage_payload
3 | 76.0 | 0A8AD70408011D6CFC6DC21DE17991C21D02EFA9C21D5A52C0C
21DD41BDAC21D4653F2C21D1F93FAC21D94EBE6C21DD5F6A9C21DA9520
BC21D4C30CD411D07A2AF421D649A0E431D3A2239431D4A1658431D1A6
56B431D068473431D308970431DB21261431D7EE045431DBA6721431D5
0EFED421DDB5096421D632E07421DD8A539C01DC5F204C21D28B15F...
```





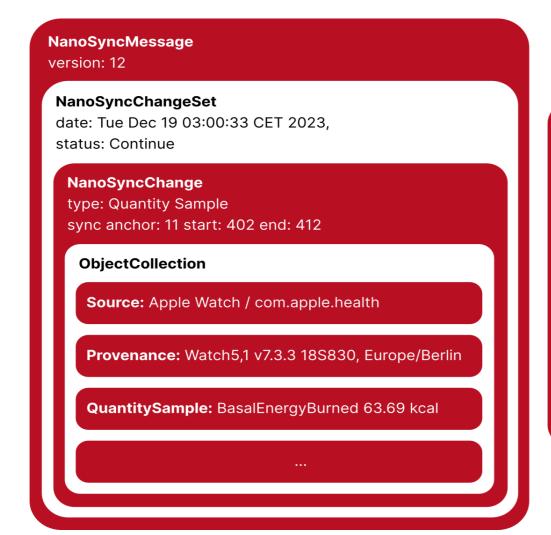
2. Within the Watch



ok cool

we started working out





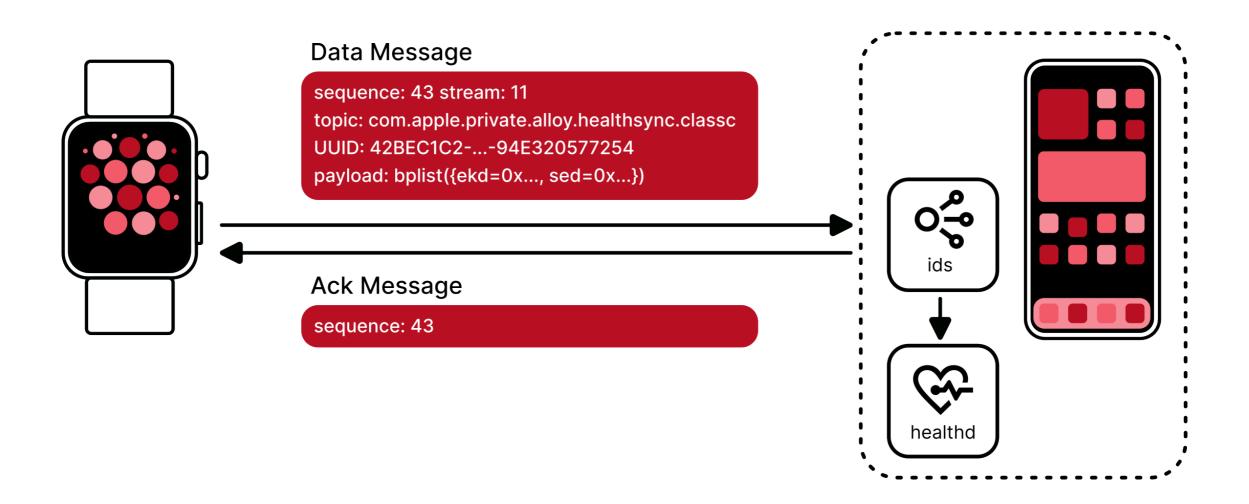


ok cool

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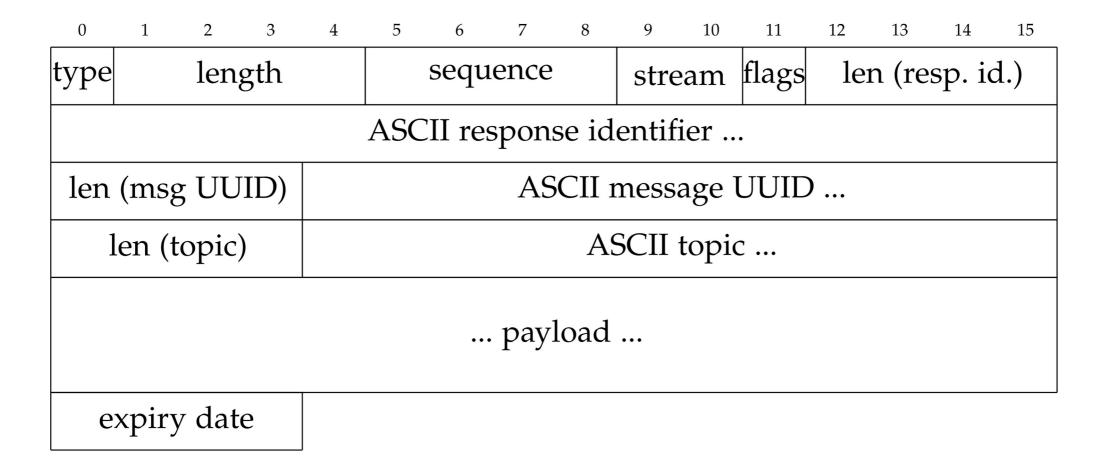


Alloy



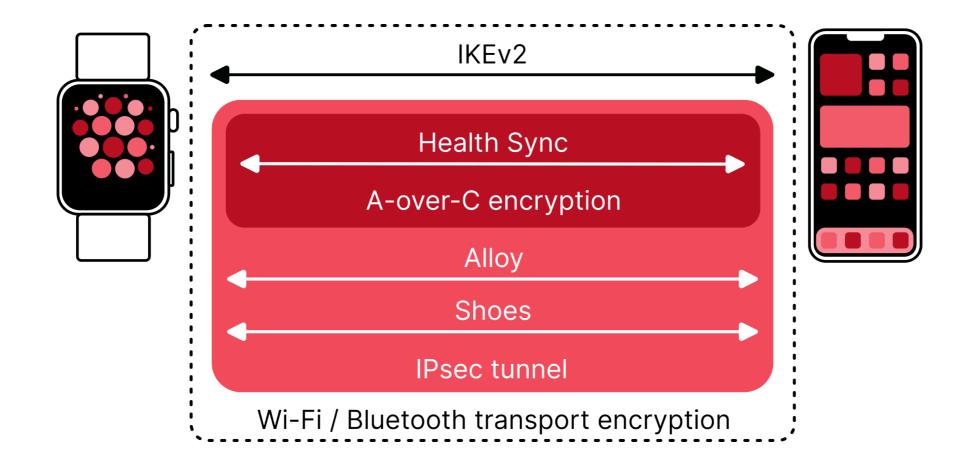


Alloy



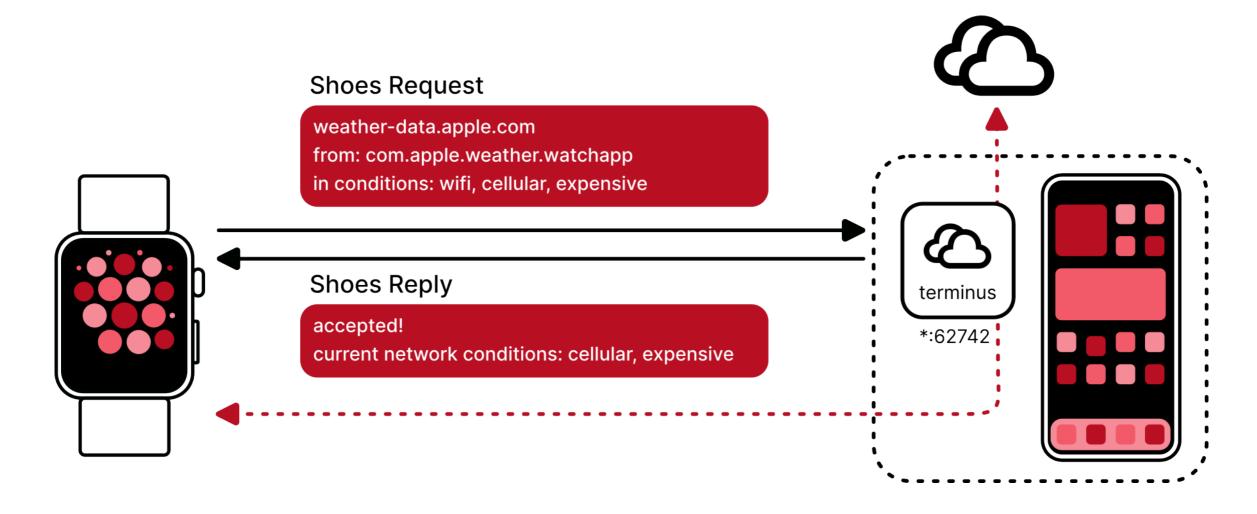


Tunnels



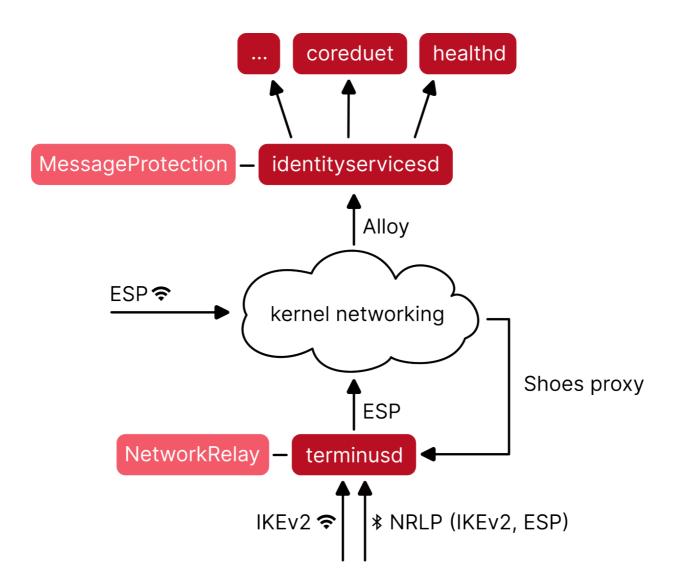


Shoes





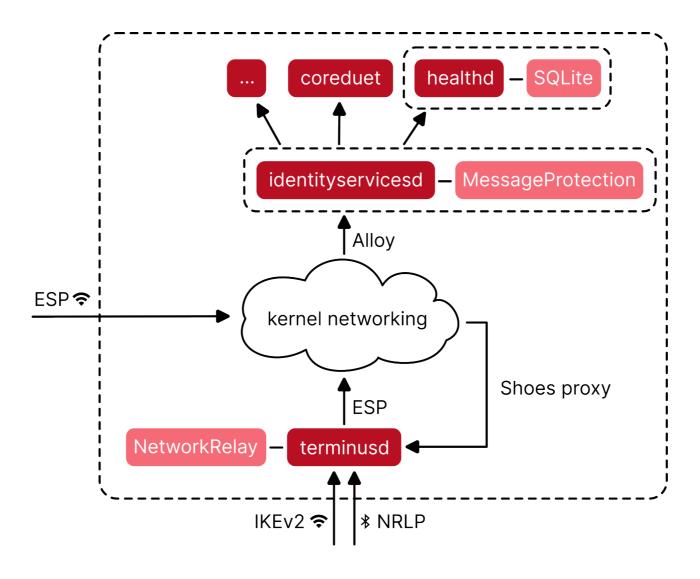
Message Flow





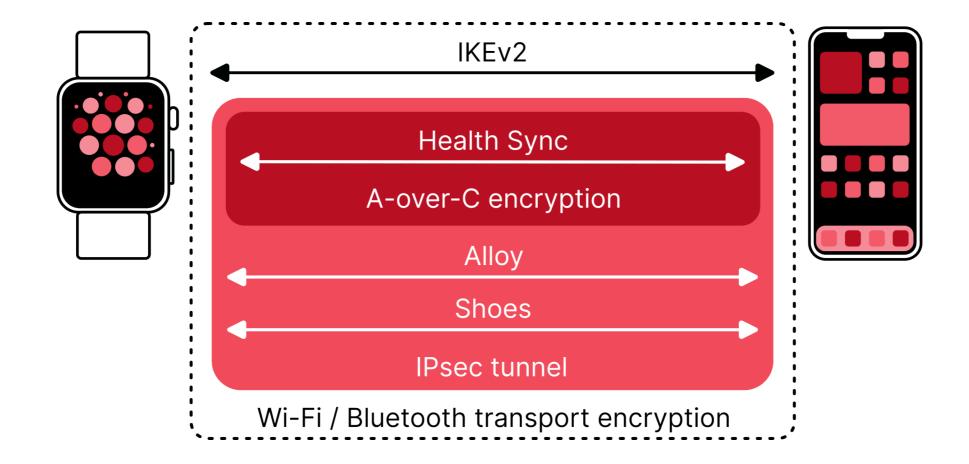


3. Attack!



trust boundaries in message handling

Re: Tunnels



IKEv2 Handling

3.10. Notify Payload

The Notify payload, denoted N in this document, is used to transmit informational data, such as error conditions and state transitions, to an IKE peer. A Notify payload may appear in a response message (usually specifying why a request was rejected), in an INFORMATIONAL exchange (to report an error not in an IKE request), or in any other message to indicate sender capabilities or to modify the meaning of the request.

— RFC 7296

48603	Device name	e.g. "iPhone", "Apple Watch"
48604	Build version	e.g. "18H17", "18S830"
50701	ProxyNotify	IPv6 address and port of SHOES server on the phone
50702	LinkDirectorMessage	used for link state signaling and WiFi discovery

Apple-defined private notify types

'	Type	Name	Comment
	1	Hello	no payload, signals restart
	2	UpdateWiFiAddressIPv6	2 byte port followed by 16 byte IP
	3	UpdateWiFiAddressIPv4	2 byte port followed by 4 byte IP
	4	UpdateWiFiSignature	variable length, unused?
	5	PreferWiFi	no payload
	6	DeviceLinkState	1 byte preferred link, 1: Bluetooth, 2: WiFi

IKEv2 Handling

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IKEv2 Handling

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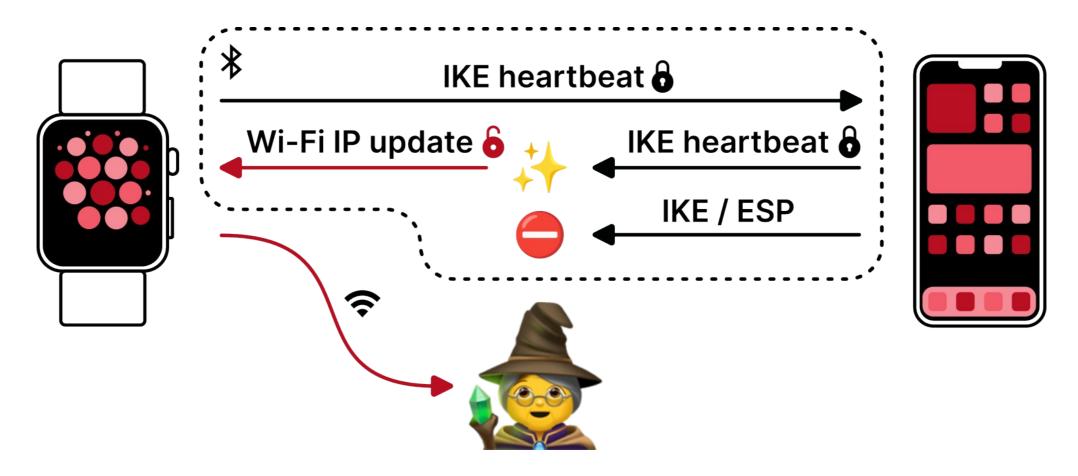
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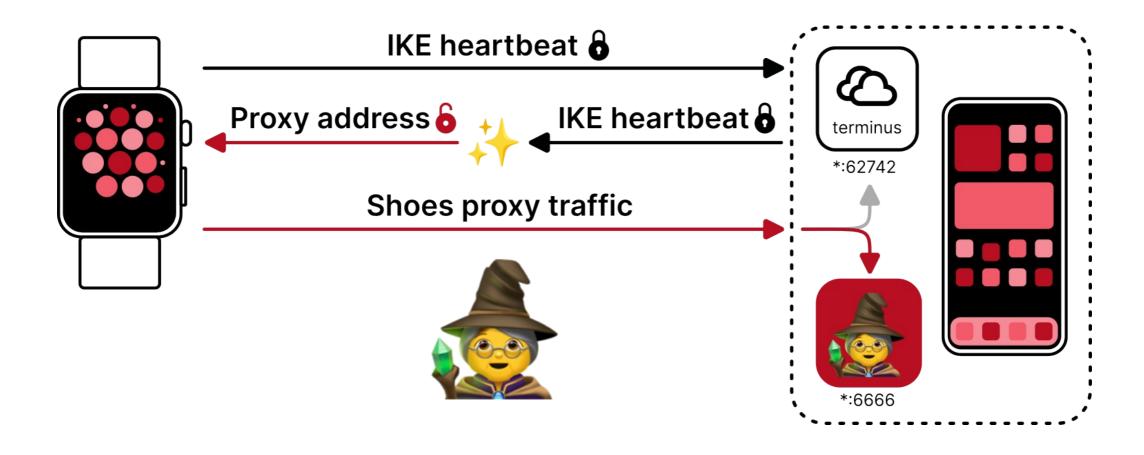
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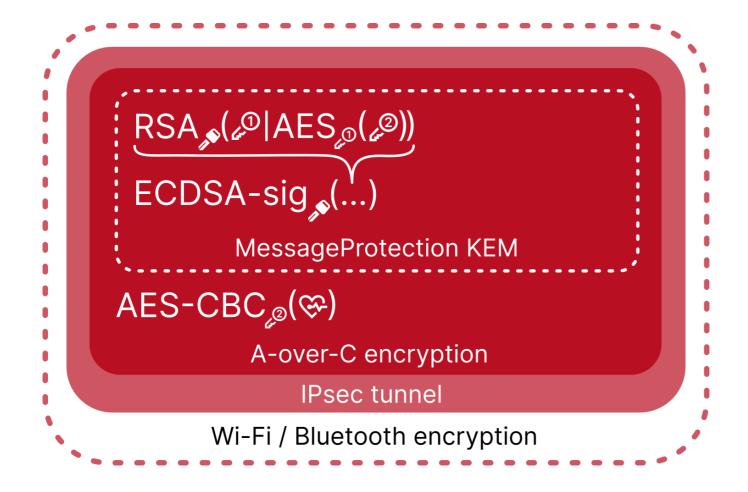
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injection of a forged Wi-Fi IP address update



injection of a forged Shoes proxy endpoint









Dancing on the Lip of the Volcano: Chosen Ciphertext Attacks on Apple iMessage

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Abstract

Apple's iMessage is one of the most widely-deployed end-to-end encrypted messaging protocols. Despite its broad deployment, the encryption protocols used by iMessage have never been subjected to rigorous cryptanalysis. In this paper, we conduct a thorough analysis of iMessage to determine the security of the protocol against a variety of attacks. Our analysis shows that iMessage has significant vulnerabilities that can be exploited by a sophisticated attacker. In particular, we outline a novel chosen ciphertext attack on Huffman compressed data, which allows retrospective decryption of some iMessage payloads in less than 218 queries. The practical implication of these attacks is that any party who gains access to iMessage ciphertexts may potentially decrypt them remotely and after the fact. We additionally describe mitigations that will prevent these attacks on the protocol, without breaking backwards compatibility. Apple has deployed our mitigations in the latest iOS and OS X releases.

1 Introduction

The past several years have seen widespread adoption of end-to-end encrypted text messaging protocols. In this work we focus on one of the most popular such protocols: Apple's iMessage. Introduced in 2011, iMessage is an end-to-end encrypted text messaging system that supports both iOS and OS X devices. While Apple does not provide up-to-date statistics on iMessage usage, in February 2016 an Apple executive noted that the system had a peak transmission rate of more then 200,000 messages per second, across 1 billion deployed devices [12].

The broad adoption of iMessage has been controversial, particularly within the law enforcement and national security communities. In 2013, the U.S. Drug Enforcement Agency deemed iMessage "a challenge for DEA intercept" [22], while in 2015 the U.S. Department of Justice accused Apple of thwarting an investigation by refusing to turn over iMessage plaintext [11]. iMessage has been at the center of a months-long debate initiated by U.S. and overseas officials over the implementation of "exceptional access" mechanisms in end-to-end encrypted communication systems [7, 26, 33], and some national ISPs have temporarily blocked the protocol [32]. Throughout this controversy, Apple has consistently maintained that iMessage encryption is end-to-end and that even Apple cannot recover the plaintext for messages transmitted through its servers [10].

Given iMessage's large installed base and the high stakes riding on its confidentiality, one might expect iMessage to have received critical attention from the research community. Surprisingly, there has been very little analysis of the system, in large part due to the fact that Apple has declined to publish the details of iMessage's encryption protocol. In this paper we aim to remedy this situation. Specifically, we attempt to answer the following question: how secure is Apple iMessage?

Our contributions. In this work we analyze the iMessage protocol and identify several weaknesses that an attacker may use to decrypt iMessages and attachments. While these flaws do not render iMessage completely insecure, some flaws reduce the level of security to that of the TLS encryption used to secure communications between enduser devices and Apple's servers. This finding is surprising given the protection claims advertised by Apple [10]. Moreover, we determine that the flaws we detect in iMessage may have implications for other aspects of Apple's ecosystem, as we discuss below.

To perform our analysis, we derived a specification for iMessage by conducting a partial black-box reverse engineering of the protocol as implemented on multiple iOS and OS X devices. Our efforts extend a high-level protocol overview published by Apple [9] and two existing partial reverse-engineering efforts [1, 34]. Armed with a protocol specification, we conducted manual cryptanal-

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Dancing on the Lip of the Volcano: Chosen Ciphertext Attacks on Apple iMessage

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Figure 2: Example of a simple ciphertext replay.

Dancing on the Lip of the Volcano: Chosen Ciphertext Attacks on Apple iMessage

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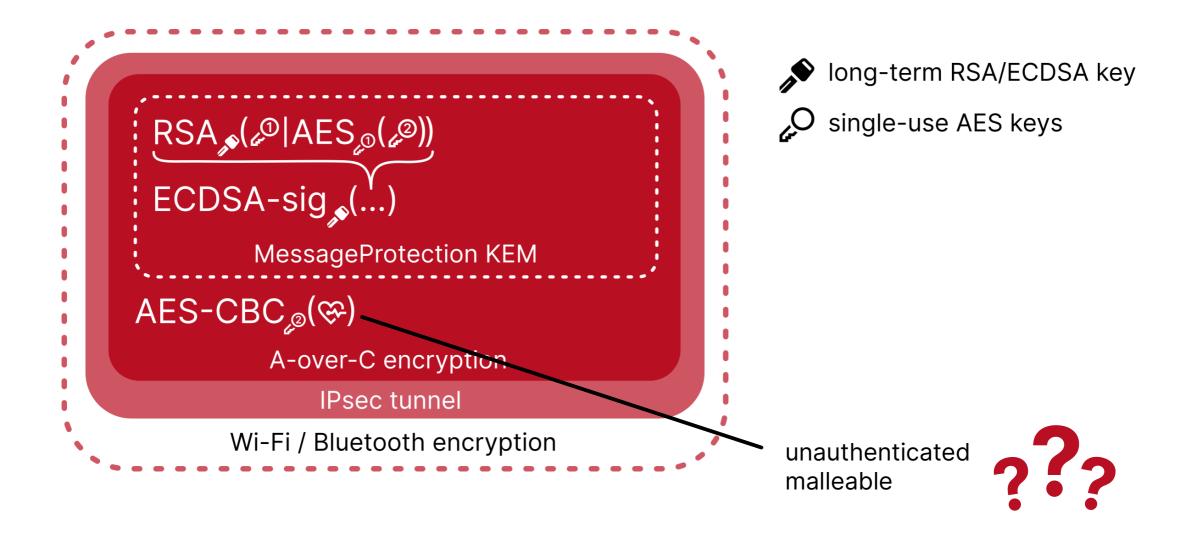
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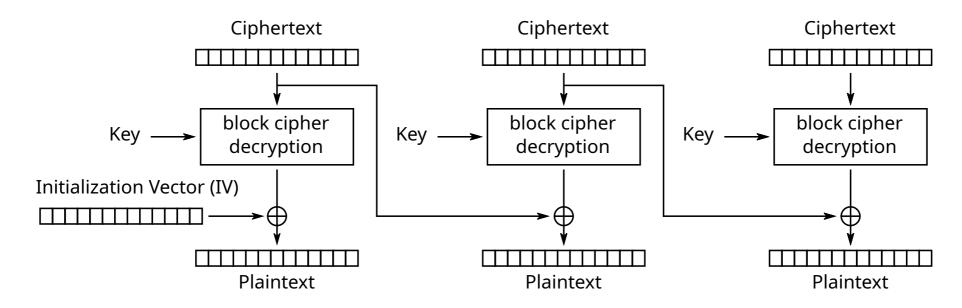
USENIX Association

- → replayability
- → malleability
- → no forward secrecy
- → compress-then-encrypt
- → weird custom crypto



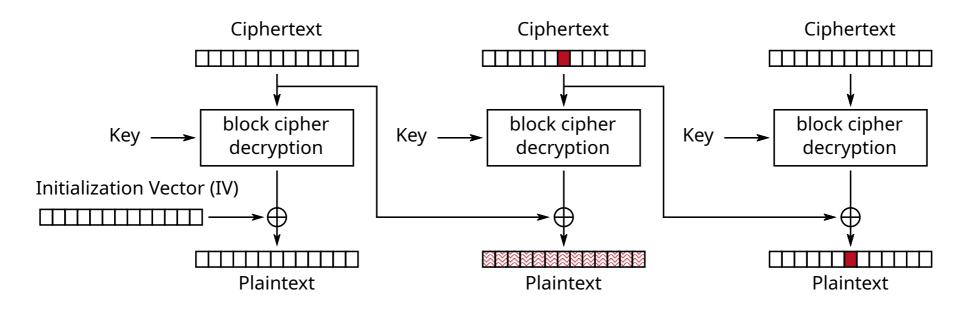






Cipher Block Chaining (CBC) mode decryption





Cipher Block Chaining (CBC) mode decryption



```
02 00 00 10 0c 1a 10 38 ce e8 1f d9 af 4a a3 af
                                             various headers
d0 68 89 a8 53 c5 41 38 00 <mark>22 9c 01 0a 83 01 0</mark>a
6d 0a 10 1e cf 3b 53 73 7c 41 64 a1 e7 f3 ee 8a
                                             heartrate sample
c5 41 19 00 00 00 00 00 80 4f 40 22 09 63 6f 7
6e 74 2f 6d 69 6e 28 01
                                             active energy sample
                                             random UUID
                                             type
                                             active energy sample
a7 73 c5 41 21 75 8b 59 2a a7 73 c5 41 11 27 3
08 ac 1c 5a c4 3f 28 01
```

genuine health sync plaintext



```
02 00 00 10 0c 1a 10 38 ce e8 1f d9 af 4a a3 af
                                                 various headers
le8 1f d9 af 4a a3 af 5f b5 59 ce 34 dc ec 31 1c
d0 68 89 a8 53 c5 41 38 00 22 9c 01 0a 83 01 0a
6d 0a 10 1e cf 3b 53 73 7c 41 64 a1 e7 f3 ee 8a
                                                 heartrate sample
c5 41 19 00 00 00 00 00 80 4f 40 22 09 63 6f 75
6e 74 2f 6d 69 6e 28 01
                                                  active energy sample
                                                 random UUID
                                                 type
                       22 3e 0a 31 0a 1b 0a 10
                                                 active energy sample
a7 73 c5 41 21 75 8b 59 2a a7 73 c5 41 11 27 31
08 ac 1c 5a c4 3f 28 01
```

04 b9 00 c2 9d e8 9a 75 98 19 0d 0c 66 cd cc fd 48 d0 4c c1 67 e7 d7 b7 <mark>85 5c 6f 7b f9 78 d</mark> 88 b5 a9 3a 61 9d e9 9c 5d 25 17 8a 7d e3 e9 73 ac 07 54 ac d2 4e f3 55 6f 2e 55 c6 91 15 6f 83 be f3 63 08 e8 67 b1 85 ee 2f e9 79 9c 3a c4 86 015cd92497ebece1

c2 23 84 52 af 1f 02 ac 7a 26 df 9b 31 d8 f1 a0

genuine health sync plaintext

A-over-C ciphertext

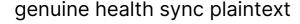


```
02 00 00 10 0c 1a 10 38 ce e8 1f d9 af 4a a3 af
                                                 various headers
e8 1f d9 af 4a a3 af 5f b5 59 ce 34 dc ec 31 1c
d0 68 89 a8 53 c5 41 38 00 22 9c 01 0a 83 01 0a
6d 0a 10 1e cf 3b 53 73 7c 41 64 a1 e7 f3 ee 8a
                                                 heartrate sample
c5 41 19 00 00 00 00 00 80 4f 40 22 09 63 6f 7!
6e 74 2f 6d 69 6e 28 01
                                                  active energy sample
                                                 random UUID
                                                 type
a7 73 c5 41 21 75 8b 59 2a a7 73 c5 41 11 27 3:
08 ac 1c 5a c4 3f 28 01
```

active energy sample

```
c2 23 84 52 af 1f 02 ac 7a 26 df 9b 31 d8 f1 a0
04 b9 00 c2 9d e8 9a 75 98 19 0d 0c 66 cd cc fo
48 d0 4c c1 67 e7 d7 b7
                             7b
 8 b5 a9 3a 61 9d e9 9c 5d 25 17 8a 7d e3 e9 7.
ac 07 54 ac d2 4e f3 55 6f 2e 55 c6 91 15 6f 83
be f3 63 08 e8 67 b1 85 ee 2f e9 79 9c 3a c4 86
015cd92497ebece1
```

A-over-C ciphertext





Malleable Encryption

```
02 00 00 10 0c 1a 10 38 ce e8 1f d9 af 4a a3 af
                                                various headers
e8 1f d9 af 4a a3 af 5f b5 59 ce 34 dc ec 31 1c
d0 68 89 a8 53 c5 41 38 00 22 9c 01 0a 83 01 0a
6d 0a 10 1e cf 3b 53 73 7c 41 64 a1 e7 f3 ee 8a
                                                heartrate sample
c5 41 19 00 00 00 00 00 80 4f 40 22 09 63 6f 75
6e 74 2f 6d 69 6e 28 01
                                                active energy sample
|654ad6214ce947140a1ed009a04660ed
                                                random UUID
                            05
                                                type
                                                active energy sample
a7 73 c5 41 21 75 8b 59 2a a7 73 c5 41 11 27 3:
08 ac 1c 5a c4 3f 28 01
```

c2 23 84 52 af 1f 02 ac 7a 26 df 9b 31 d8 f1 a0 04 b9 00 c2 9d e8 9a 75 98 19 0d 0c 66 cd cc fo 48 d0 4c c1 67 e7 d7 b7 7b be f3 63 08 e8 67 b1 85 ee 2f e9 79 9c 3a c4 86 015cd92497ebece1

genuine health sync plaintext

A-over-C ciphertext



Malleable Encryption

```
02 00 00 10 0c 1a 10 38 ce e8 1f d9 af 4a a3 af
                                                 various headers
e8 1f d9 af 4a a3 af 5f b5 59 ce 34 dc ec 31 1c
d0 68 89 a8 53 c5 41 38 00 <mark>22 9c 01 0a 83 01 0</mark>a
6d 0a 10 1e cf 3b 53 73 7c 41 64 a1 e7 f3 ee 8a
                                                 heartrate sample
c5 41 19 00 00 00 00 00 80 4f 40 22 09 63 6f 7!
6e 74 2f 6d 69 6e 28 01
                                                  heart rate sample
|654ad6214ce947140a1ed009a04660ed
                                                 random UUID
                             05
                                                 type
                                                 active energy sample
a7 73 c5 41 21 75 8b 59 2a a7 73 c5 41 11 27 3:
08 ac 1c 5a c4 3f 28 01
```

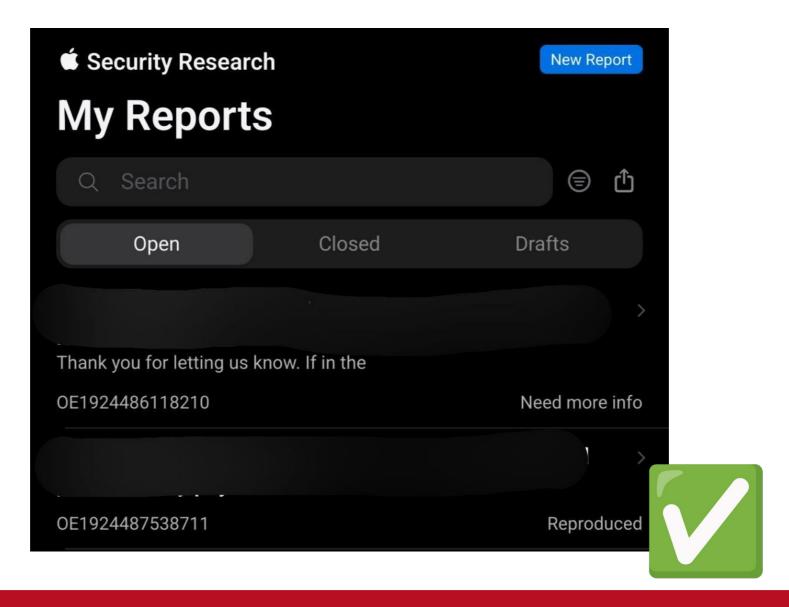
```
c2 23 84 52 af 1f 02 ac 7a 26 df 9b 31 d8 f1 a0
04 b9 00 c2 9d e8 9a 75 98 19 0d 0c 66 cd cc fo
48 d0 4c c1 67 e7 d7 b7
                       7b
be f3 63 08 e8 67 b1 85 ee 2f e9 79 9c 3a c4 86
015cd92497ebece1
```

A-over-C ciphertext

genuine health sync plaintext



Responsible Disclosure





Security Takeaways

- Standards exist for a reason
- Crypto will (not) save you
- Consider unexpected system interactions
- Think really really hard before rolling your own crypto
- Avoid complexity whereever possible



SMSDownloadOutgoing GenericCommandMessage DataMessage AckMessage SMSDeliveryReceipt GenericGroupMessageCommand SMSReadReceipt LocationShareOfferCommand KeepAliveMessage ProtobufMessage ExpiredAckMessage **SMSFailure** Handshake FragmentedMessage ErrorMessage **EncryptedMessage** ResourceTransferMessage ServiceMapMessage DictionaryMessage **OTREncryptedMessage** SessionReinitiateMessage **AppAckMessage OTRMessage** SyndicationAction ProxyOutgoingNiceMessage SessionInvitationMessage RetractMessage ProxyIncomingNiceMessage SessionAcceptMessage EditMessage SessionDeclineMessage TextMessage RecoverSyncMessage SessionCancelMessage DeliveryReceipt MarkAsUnreadMessage DeliveredQuietlyMessage SessionMessage ReadReceipt NotifyRecipientMessage SessionEndMessage AttachmentMessage SMSTextMessage PlayedReceipt RecoverJunkMessage SMSFilteringSettingsMessage SMSTextDownloadMessage SavedReceipt

ReflectedDeliveryReceipt



SMSOutgoing

accessibility.local accessibility.switchcontrol accounts.representative accountssync addressbooksync airtr timers amsaccountsync anisette appconduit appconduit.v2 applepay applepay.identitycredential applepay.shar apppredictionsync appregistrysync appstore appsyncconduit appsyncconduit.v2 arcade.fastsync askto audiocon audiocontrol.music autobugcapture avconference.avctester biz bluetooth.audio bluetoothregistry bluetoothre bluetoothregistryclassa bluetoothregistryclassc brook bulletinboard bulletindistributor bulletindistributo callhistorysync camera.proxy carmelsync carousel.uitrigger clockface.sync cmsession companion.auth compani contextsync contextsync.local continuity.activity continuity.auth continuity.auth.classa continuity.encryp continuity.tethering continuity.unlock coreduct coreduct.sync ct.baseband.p2p.notification ct.commcenter.p ct.commcenter.sim ct.commcenter.sim.cloud ded ded.watch digitalhealth donotdisturb dropin.communication dr electrictouch eventkitmutation eventkitsync facetime.audio facetime.lp facetime.messaging facetime.multi f facetime.sync facetime.video familycontrols fignero findmy.itemsharing-crossaccount findmydeviced.aoverc findmydeviced.watch fi fmd fmd.local fmf fmf. gamecenter gelato gfta ol.cloud groupRemoteCo groupRemoteControl.ses haringsetup healthapp. healthappnotifications emoteurlconnection ids ~230 Alloy topics across 151 iOS binaries intercom internal.watd manager itunes itunesc kbd.transfer kcsharind .auth location.fenceha

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4. Enter the Witch

Case 2:24-cv-04055 Document 1 Filed 03/21/24 Page 1 of 88 PageID: 1

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IN THE UNITED STATES DISTRICT COURT FOR THE DISTRICT OF NEW JERSEY

UNITED STATES OF AMERICA

U.S. Department of Justice, Antitrust Division 450 Fifth Street NW, Suite 8600 Washington, DC 20530

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STATE OF MAINE

6 State House Station

No.



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Case 2:24-cv-04055 Document 1 Filed 03/21/24 Page 12 of 88 PageID: 12

to set a stake in the ground for what features we think are 'good enough' for the consumer. I would argue we're already doing *more* than what would have been good enough. But we find it very hard to regress our product features YOY [year over year]." Existing features "would have been good enough today if we hadn't introduced [them] already," and "anything new and especially expensive needs to be rigorously challenged before it's allowed into the consumer phone." Thus, it is not surprising that Apple spent more than twice as much on stock buybacks and dividends as it did on research and development.

- 15. Moreover, Apple has demonstrated its ability to use its smartphone monopoly to impose fee structures and manipulate app review to inhibit aggrieved parties from taking advantage of regulatory and judicial solutions imposed on Apple that attempt to narrowly remedy harm from its conduct.
- 16. Apple wraps itself in a cloak of privacy, security, and consumer preferences to justify its anticompetitive conduct. Indeed, it spends billions on marketing and branding to promote the self-serving premise that only Apple can safeguard consumers' privacy and security interests. Apple selectively compromises privacy and security interests when doing so is in Apple's own financial interest—such as degrading the security of text messages, offering governments and certain companies the chance to access more private and secure versions of app stores, or accepting billions of dollars each year for choosing Google as its default search engine when more private options are available. In the end, Apple deploys privacy and security justifications as an elastic shield that can stretch or contract to serve Apple's financial and business interests.
- 17. Smartphones have so revolutionized American life that it can be hard to imagine a world beyond the one that Apple, a self-interested monopolist, deems "good enough." But under

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- Smartphones have
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Case 2:24-cv-04055 Document 1 Filed 03/21/24 Page 40 of 88 PageID: 40

user to purchase a different kind of smartphone because doing so requires the user to abandon their costly Apple Watch and purchase a new, Android-compatible smartwatch.

- 97. By contrast, cross-platform smartwatches can reduce iPhone users' dependence on Apple's proprietary hardware and software. If a user purchases a third-party smartwatch that is compatible with the iPhone and other smartphones, they can switch from the iPhone to another smartphone (or vice versa) by simply downloading the companion app on their new phone and connecting to their smartwatch via Bluetooth. Moreover, as users interact with a smartwatch, e.g., by accessing apps from their smartwatch instead of their smartphone, users rely less on a smartphone's proprietary software and more on the smartwatch itself. This also makes it easier for users to switch from an iPhone to a different smartphone.
- Apple recognizes that driving users to purchase an Apple Watch, rather than a third-party cross-platform smartwatch, helps drive iPhone sales and reinforce the moat around its smartphone monopoly. For example, in a 2019 email the Vice President of Product Marketing for Apple Watch acknowledged that Apple Watch "may help prevent iPhone customers from switching." Surveys have reached similar conclusions: many users say the other devices linked to their iPhone are the reason they do not switch to Android.
- 99. Apple also recognizes that making Apple Watch compatible with Android would "remove[an] iPhone differentiator."
- 100. Apple uses its control of the iPhone, including its technical and contractual control of critical APIs, to degrade the functionality of third-party cross-platform smartwatches in at least three significant ways: First, Apple deprives iPhone users with third-party smartwatches of the ability to respond to notifications. Second, Apple inhibits third-party smartwatches from maintaining a reliable connection with the iPhone. And third, Apple

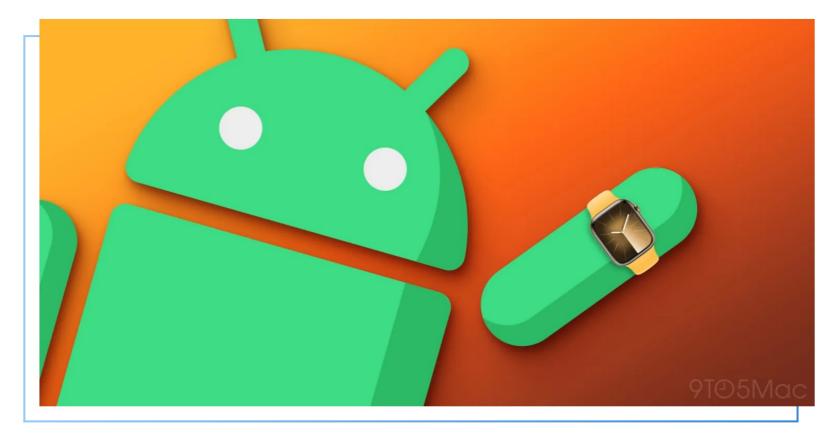
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Apple says it spent three years trying to bring Apple Watch to Android

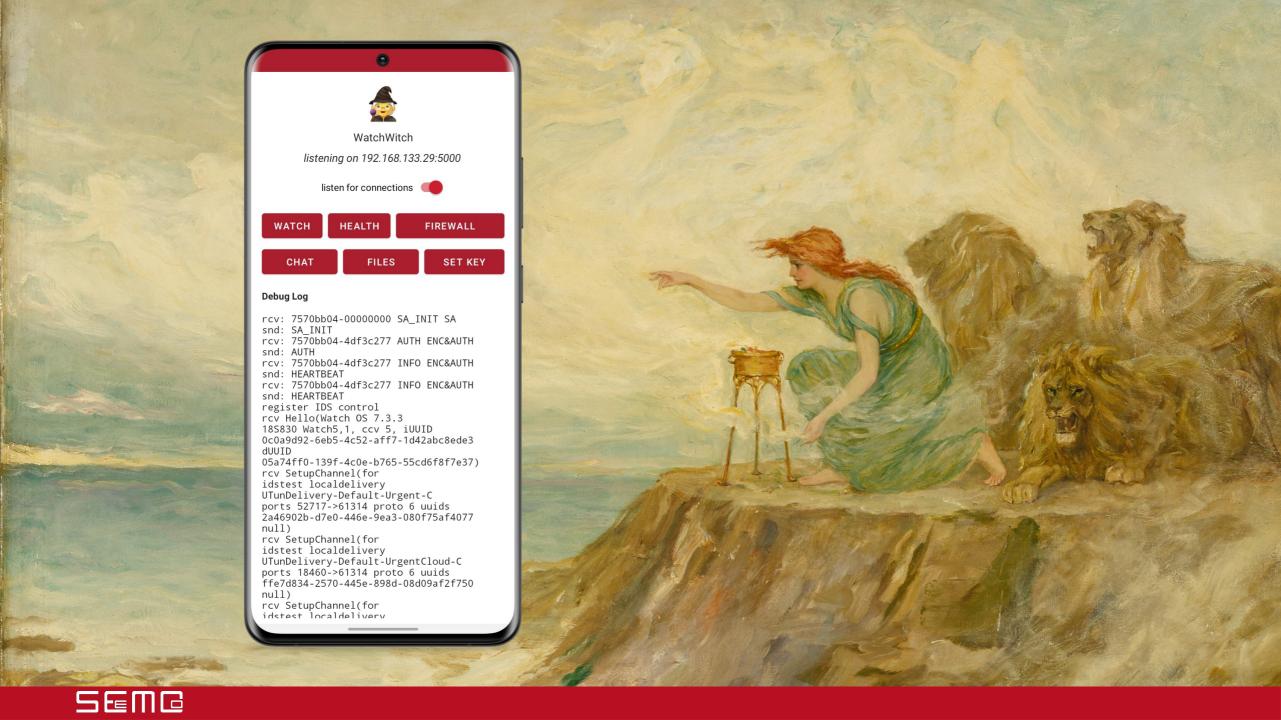


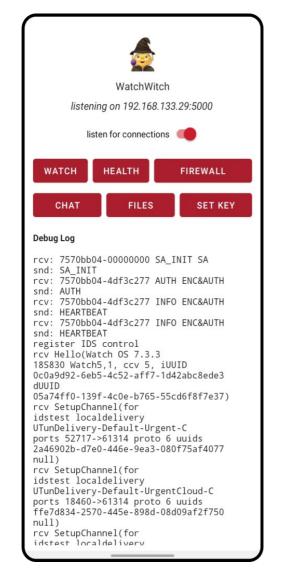
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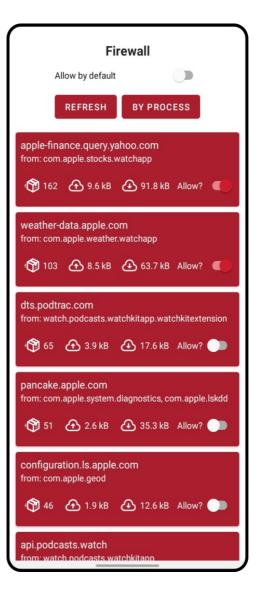
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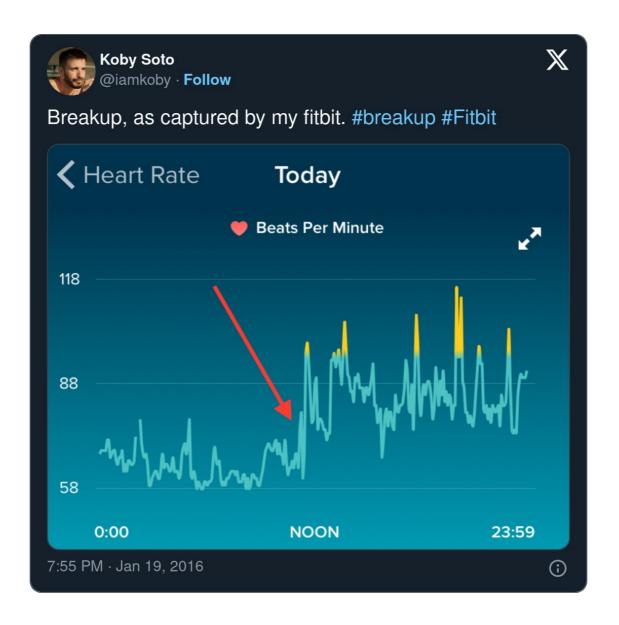






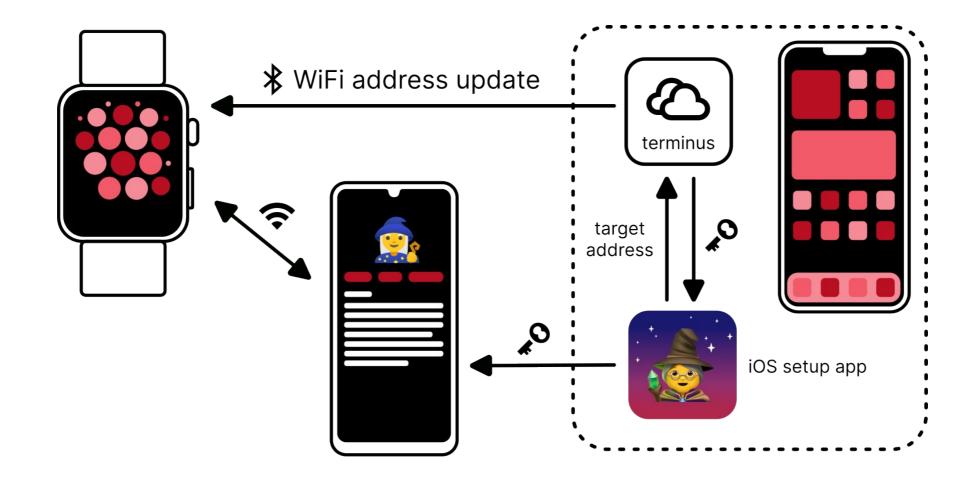






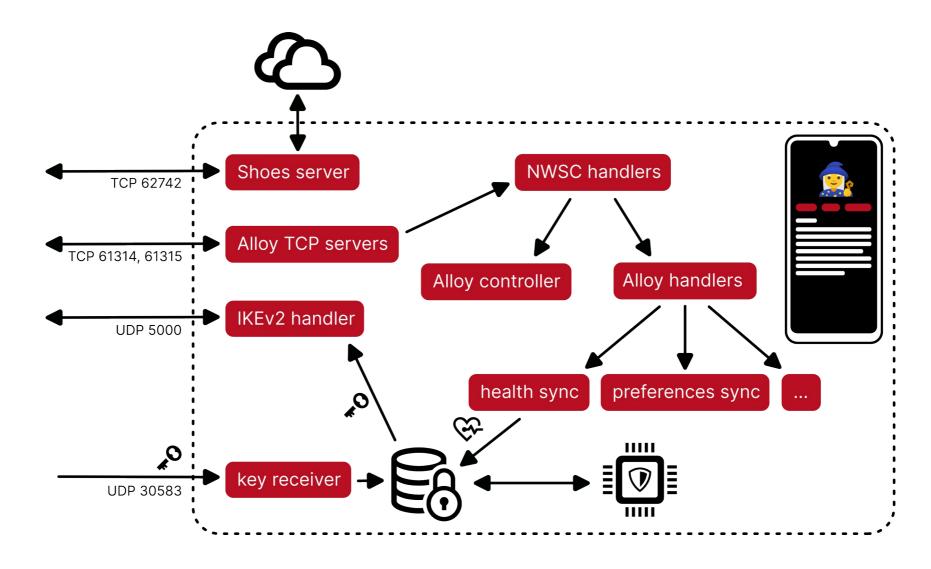


App Architecture





App Architecture



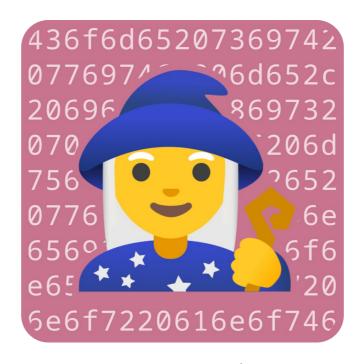
Interoperability Takeaways

- It can be done.
- It can be secure.
- Open Interfaces are curb cuts
- A better world is possible 🔆

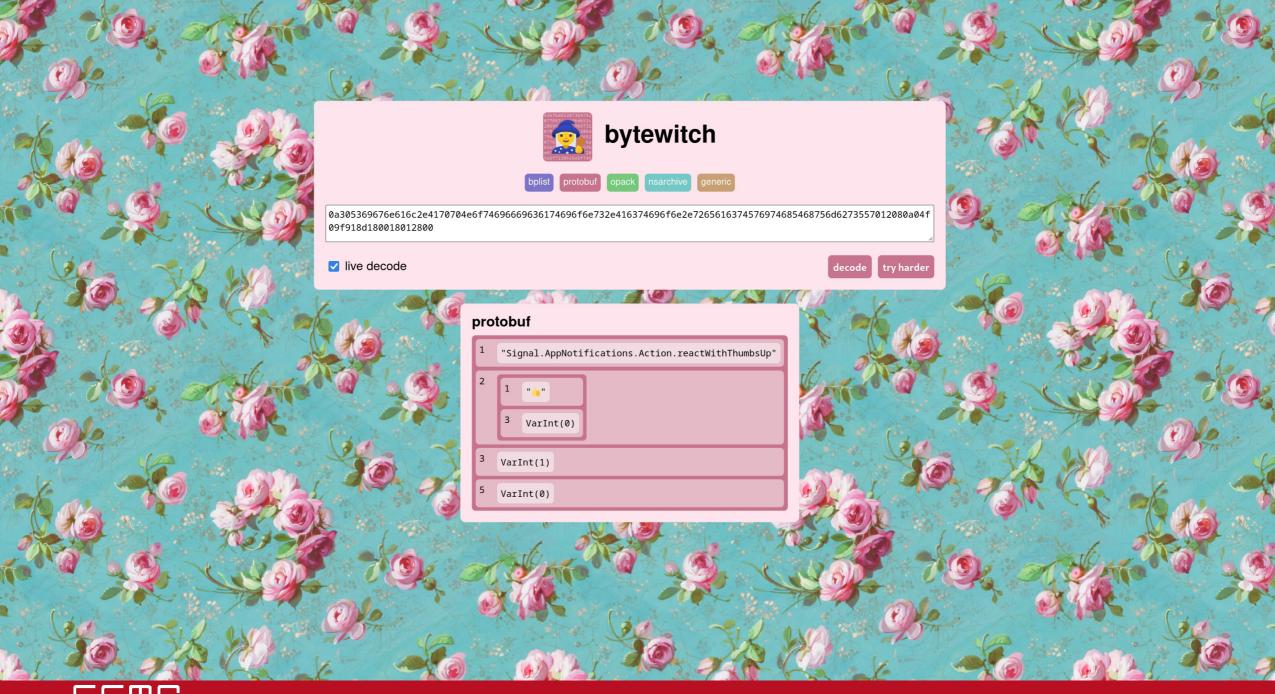


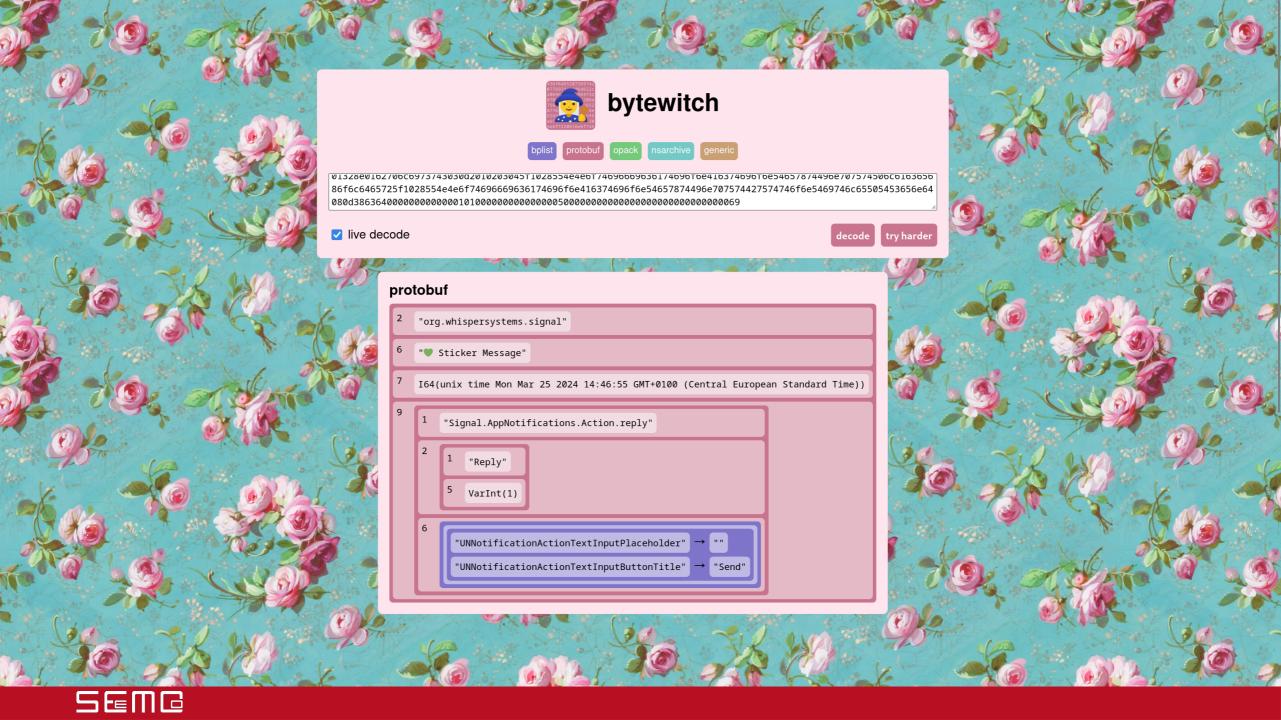


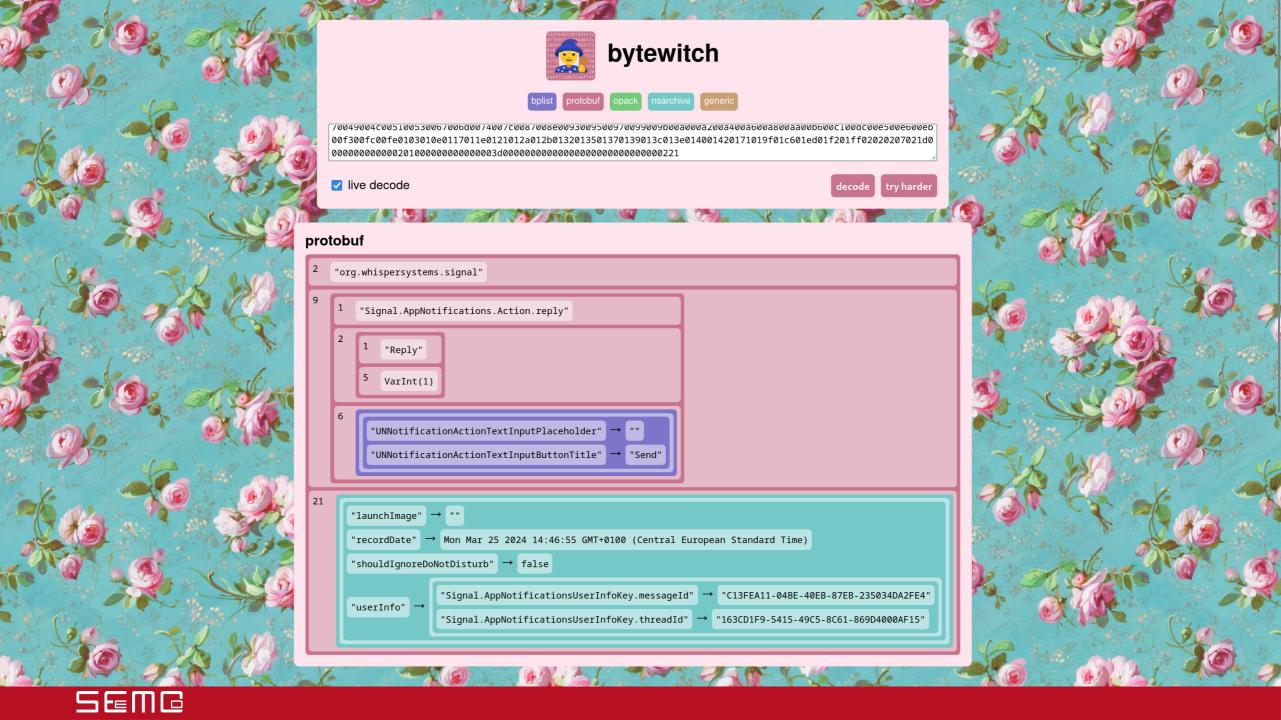
app release & source code tooling, frida scripts, wireshark dissectors more protocol documentation stay tuned **

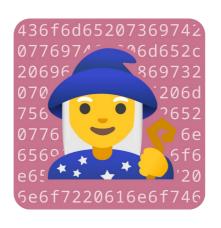


5. One more thing...









bytewitch

rec0de.net/open/bytewitch/

github.com/rec0de/bytewitch/







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