

Neohapsis is now part of Cisco.



Multipath TCP



Breaking Today's Networks
With Tomorrow's Protocols



Speakers - Who are we?

- Catherine (Kate) Pearce
 - Security Consultant / Pentester
 - New Zealand transplant.
 - Loves her wine the way she likes her RFCs (Dry)



- Patrick Thomas
 - Senior Security Consultant / Pentester
 - Application Security focus





MPTCP changes fundamental assumptions about

how TCP works*

Use it to break things today

Adapt to it for tomorrow





Not Layer 4?
Totally the same.

Layer 4? Buckle Up.



```
root@deb7min2:~# curl 192.168.88.164
<html><body><hl>is the default web page for this server.
The web server software is running but no content has been </body></html>
root@deb7min2:~#
```



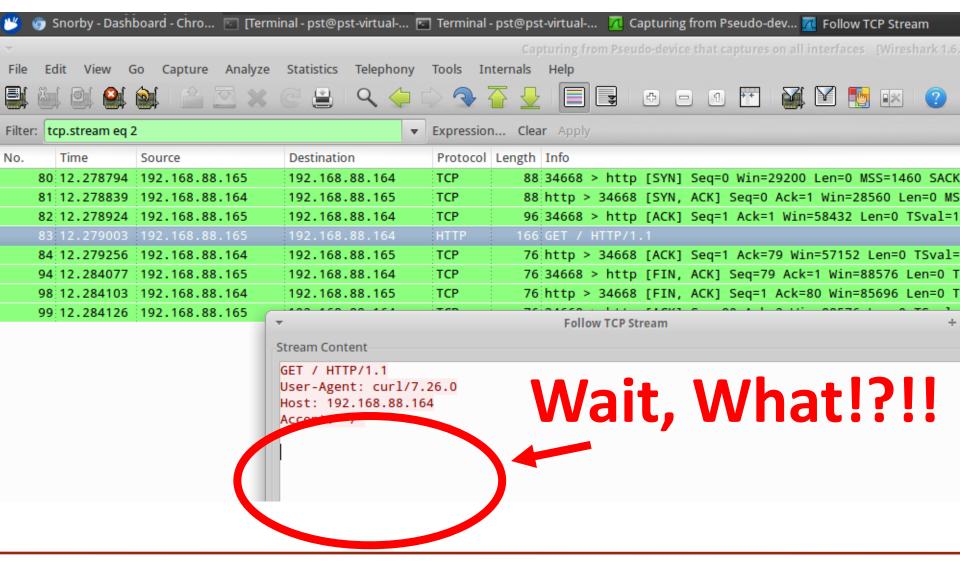


🖐 🧑 Snorby - Dashboard - Chro 📧 [Terminal - pst@pst-virtual 📧 Terminal - pst@pst-virtual 📶 Capturing from Pseudo-dev								
₩	Capturing from Pseudo-device that captures on all interfaces [Wireshark 1.6.7]							
File	File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help							
				~ @			nlar O	
			₹ @ 🚇 Q 🖕	7			* ?	
Filter: tcp.port == 80								
No.	Time	Source	Destination	Protocol	l Length Info			
	30 12.27879	94 192.168.88.165	192.168.88.164	TCP	88 34668 > http [SY	N] Seq=0 Win=29200 Len=0 MS	SS=1460 SACK_PERM=1 TSval=1249	
8	31 12.27883	39 192.168.88.164	192.168.88.165	TCP			560 Len=0 MSS=1460 SACK_PERM=1	
		04 100	192.168.88.164	TCP	96:34668 > http://doi.org/10.000	.K] Seq-1 Ack	en=0 TSval=12490113 TSecr=1247	
	83 1 7900	3 192.168.88.165	192.168.88.164	HTTP	166 GET / HTTP/1.1	Mark Packet (toggle)		
8	84 12.27925	0 192.100	100 160 00 165	TCP	76:http > 24669 FAC	Ignore Packet (toggle)	0 TSval=12474351 TSecr=124	
8	35 12.28009	95 192.168.88.165	192.168.88.164	TCP	88 39757 > http [SY	Cot Time Reference (toggle)	460 SACK_PERM=1 TSval=1249	
8	36 12.28011	11 192.168.88.164	192.168.88.165	TCP	92 http > 39757 [SY	Set Time Reference (toggle)	Len=0 MSS=1460 SACK_PERM=1	
8	37 12.28022	192.168.88.165	192.168.88.164	TCP	92 39757 > http [AC	Manually Resolve Address	TSval=12490113 TSecr=1247	
1	38 12.28033	38 192.168.88.164	192.168.88.165	TCP	76 [TCP Window Upda	Apply as Filter	ck=1 Win=85696 Len=0 TSval	
8	39 12.28398	34 192.168.88.164	192.168.88.165	HTTP	548 HTTP/1.1 200 OK	Prenare a Filter		
		14 192.168.88.165	192.168.88.164	TCP	76 39757 > http [AC	Conversation Filter	=0 TSval=12490114 TSecr=12	
9	1 12.28404	15 192.168.88.165	192.168.88.164	TCP	88 [TCP Dup ACK 90#		k=461 Win=88576 Len=0 TSva	
		192.168.88.164	192.168.88.165	TCP	88 [TCP Dup ACK 89#	Colorize Conversation	Ack=1 Win=85696 Len=0 TSva	
		73 192.168.88.165	192.168.88.164	TCP	76 39757 > b+ [F1		en=0 TSval=12490114 TSe	
		77 192.168.88.165	192.168.88.164	TCP	76 34668 > ttp [FI		=0 TSval=12490114 TSec	
		36 192.168.88.165	192.168.88.164	TCP	76 [TCP Dup ACK 95%		k=461 Win=88576 Len=0 TSva	
		91 192.168.88.164	192.168.88.165	TCP	76 http > 39757 [FI		6 Len=0 TSval=12474352 TSe	
		99 192.168.88.165	192.168.88.164	TCP	76 39757 > http [AC	СОРУ	=0 TSval=12490114 TSecr=12	
		192.168.88.164	192.168.88.165	TCP	76 http > 34668 [FI	0.0	Len=0 TSval=12474352 TSec	
9	9 12.28412	26 192.168.88.165	192.168.88.164	TCP	76 34668 > http [AC		0 TSval=12490114 TSecr=124	
						Print		
						Show Packet in New Window		
N Frame 92: 166 butes on wire (1239 bits) 166 butes contured (1239 bits)								

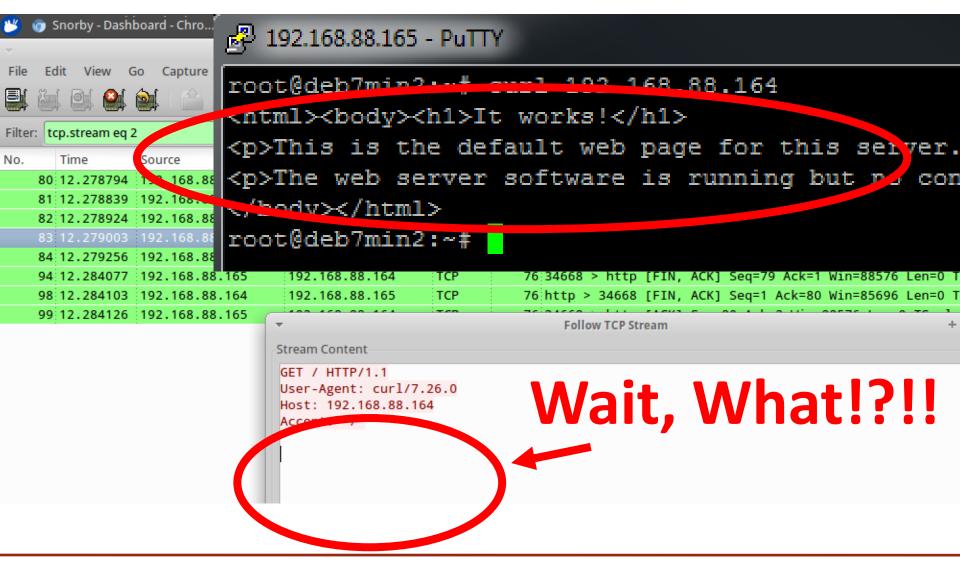
- ▶ Frame 83: 166 bytes on wire (1328 bits), 166 bytes captured (1328 bits)
- ▶ Linux cooked capture

Mynertext Transfer Protocol

- ▶ Internet Protocol Version 4, Src: 192.168.88.165 (192.168.88.165), Dst: 192.168.88.164 (192.168.88.164)
- ▶ Transmission Control Protocol, Src Port: 34668 (34668), Dst Port: http (80), Seq: 1, Ack: 1, Len: 78









nc 192.168.1.25 3000



```
root@deb7min–LEFT:/home/username# netstat
       Internet connections (w/o servers)
Proto Recv–Q Send–Q Local Address
                                               Foreign Address
                                                                         State
tcp
                                               192.168.1.25:3000
                                                                           YN_SENT
tcp
tcp
                          .3.111:33145
                                               192.168.22.145:3000
                                                                             SENT
                                                                         SY
                                               192.168.1.25:3000
                                                                              SENT
tcp
                                               192.168.22.145:3000
                                                                             SENT
                          .1.111:52605
tcp
                     192.168.1.34:5081
                                               192.168.1.25:3000
                                                                         EST
                                                                             ABLISHED
tcp
                                                                         SYN
                                                                             SENT
                     192.168.1.34:3409
                                               192.168.22.145:3000
tcp
tcp
                                                                         SYN
                                                                             SENT
                     10.3.3.111:36916
                                               192.168.1.25:3000
                                               192.168.22.145:3000
                                                                             SENT
tcp6
                     2601:6:1700:168:2:4
                                          0378
                                               2601:6:1700:168:20:3000
tcp6
                                               2601:6:1700:168:20:3000
                     2601:6:1700:168:2:41
                                                                          STABLISHED
```







Sense

This makes none



Why did we see that?

→ Let's talk about MPTCP

...but first, why change TCP?





Current TCP is rather limited

Doesn't support use cases for:

- High Availability
- Link Aggregation
- Multihoming
- Mesh networking





Multipath TCP is an extension to TCP that adds the above functionality

AND: it works over existing infrastructure

(it *IS* TCP... just more so)

BUT: nothing much else understands it

including security tools





MPTCP is more culture shock than security vulnerability

We like MPTCP
We want MPTCP to succeed
Network security isn't ready





Background **Technical Introduction Key Security Effects Perimeter Security Network Management MPTCP** Future





- Kate saw this...
- Y Hacker News new | comments | ask | jobs | submit login
- Apple seems to also believe in Multipath TCP (uclouvain.be) 10 points by rhapsodyv 278 days ago | discuss | save to pocket





Which led to this...

Hacker News new | comments | ask | jobs | submit login

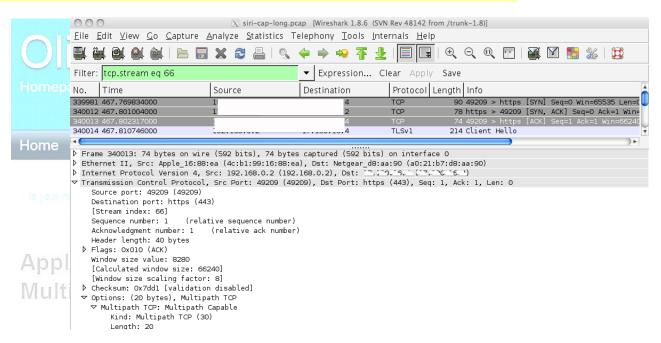
Olivier Bonaventure Homepage and blog Home About CV **Publications** Teaching People « Is your network ready for iOS7 and Multipath TCP? Quickly producing time-sequence diagrams » September 18, 2013 Apple seems to also believe in **Multipath TCP**





Which contained this...

Hacker News new | comments | ask | jobs | submit login







Then other media outlets started covering it...

Hacker News new | comments | ask | jobs | submit | login

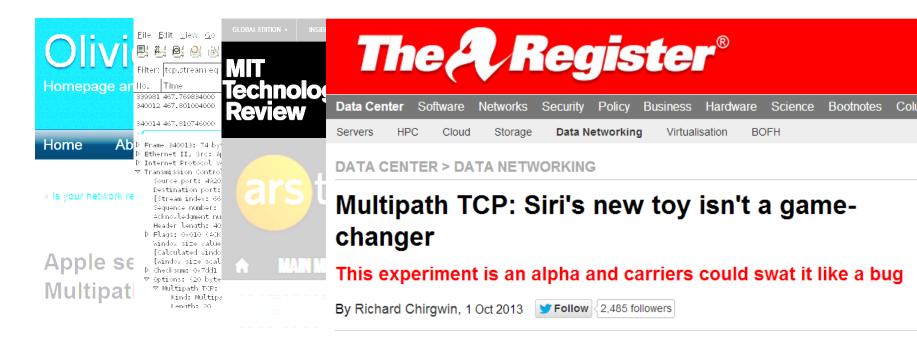






Then other media outlets started covering it...not always positively

Hacker News new | comments | ask | jobs | submit login







And then...

SILENCE

BUT, the rate of progress was unprecedented for a major change to TCP





Was anyone thinking about security?

The security of MPTCP itself



What changes like this could mean for network security



... not so much





That's what this session is about

What does multipath TCP mean for security today?

What could it (or similar tech) mean to network security a decade from now?

With a couple of PoCs and tools...





Background **Technical Introduction Key Security Effects Perimeter Security Network Management MPTCP** Future





Motivations and Advantages

TCP implements connections between IP:PORT & IP:PORT

NOT between endpoint A and endpoint B

In the past this was a distinction without a difference, but not any more





An MPTCP Connection is defined by a connection ID

It is composed of multiple streams, where each stream is a regular TCP connection (with an option strapped on)



MPTCP Characteristics

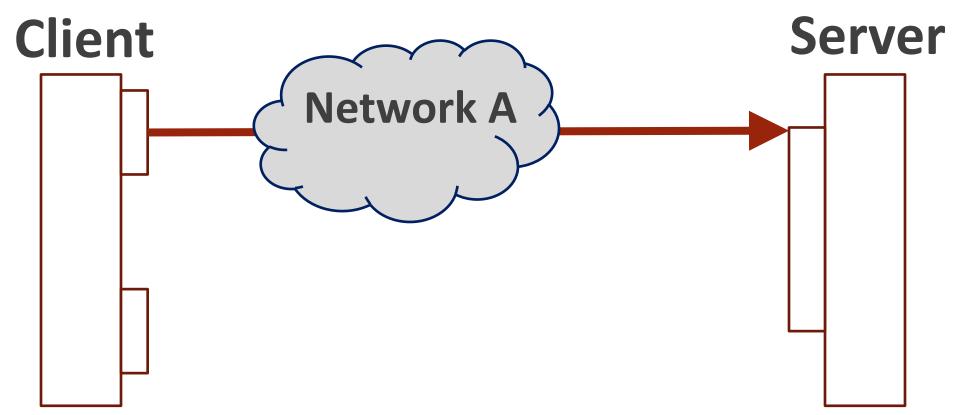
Backwards compatibility

Performance >= now

Security >= now



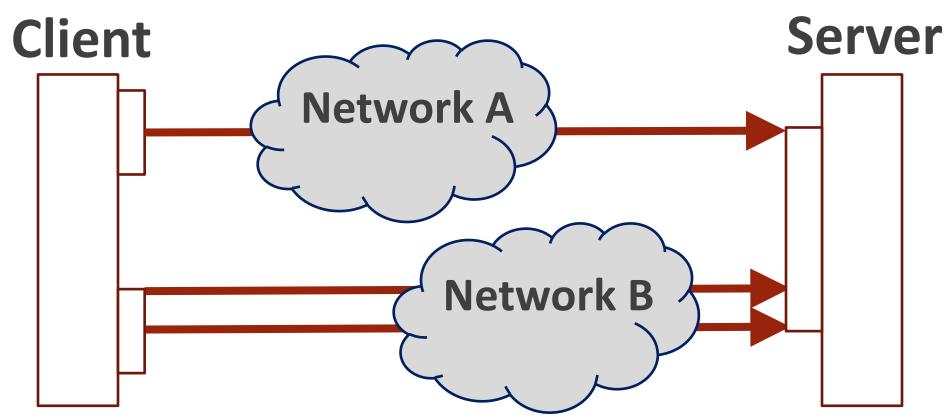




MPTCP connection looks like TCP so far...



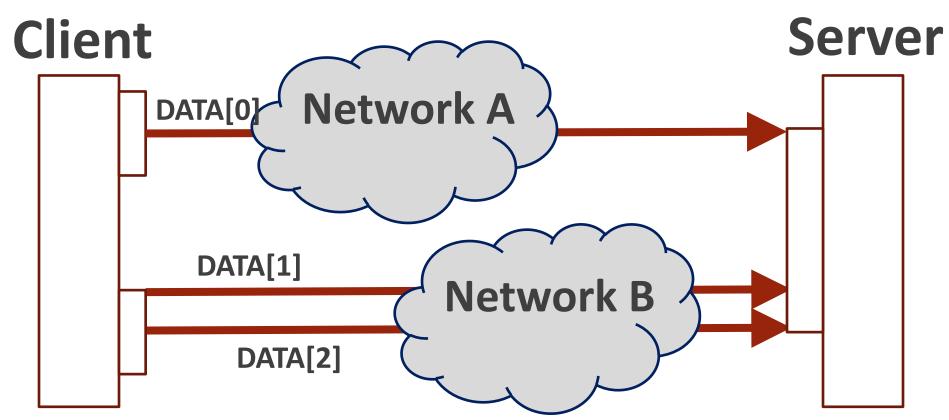




N different TCP connections, contributing to *ONE* logical data flow



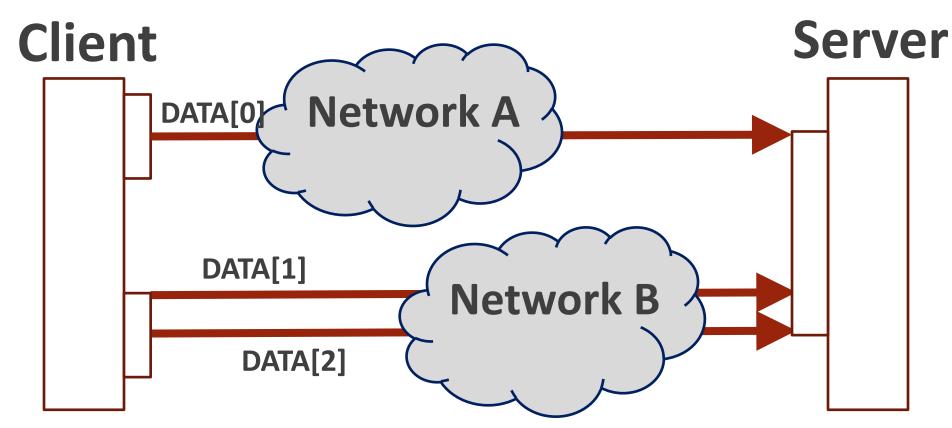
MPTCP – Simple Case



N different TCP connections, contributing to *ONE* logical data flow... data flows through any/all



MPTCP – Simple Case

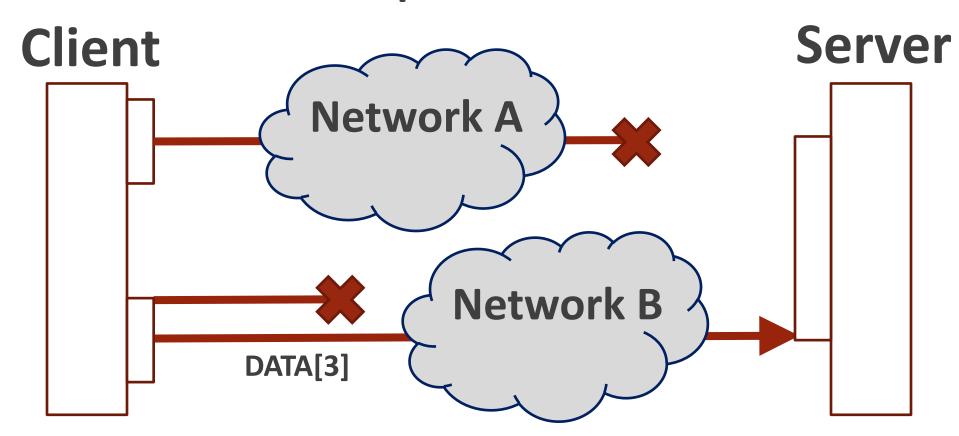


Sender of a packet can choose to use any flow *(this will be important)*





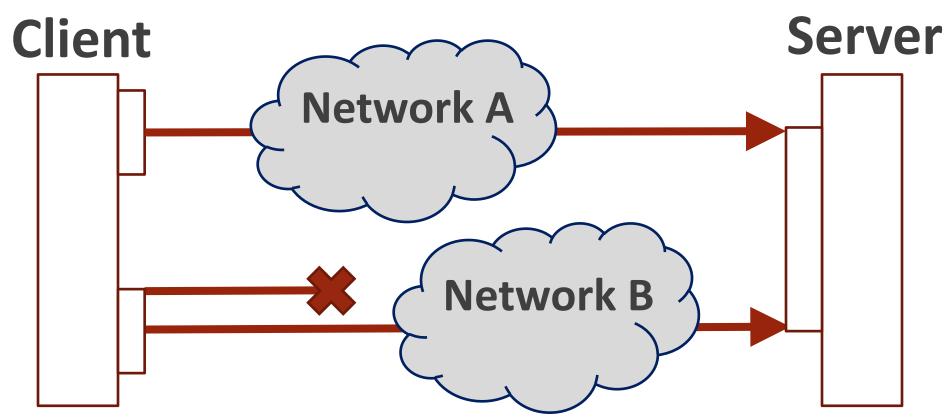
MPTCP – Simple Case



Any subset of connections can drop, overall flow continues.



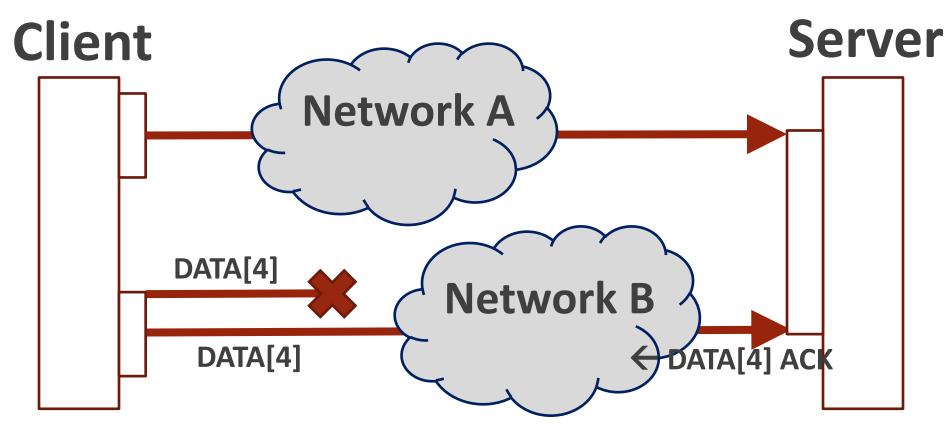




Connections can be re-added at any time



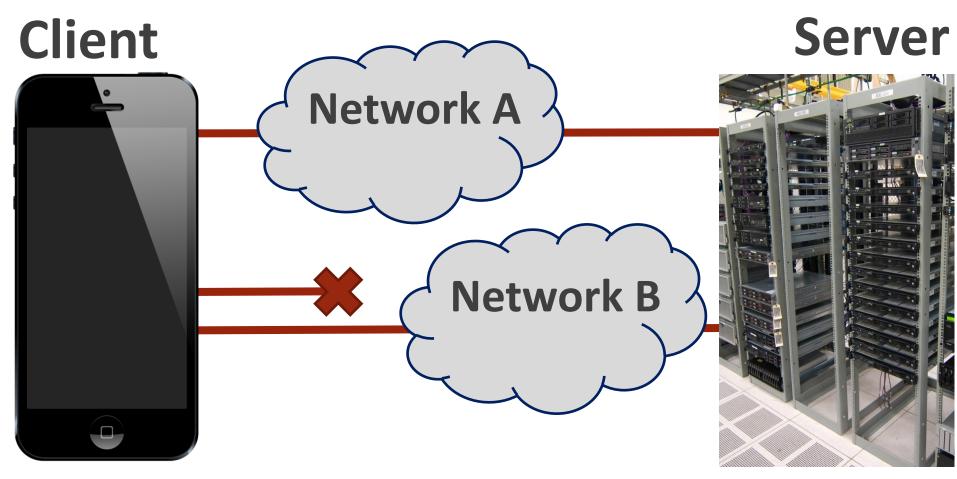




Un-ACK'd data can be quickly resent over a different flow... first ACK is good enough!



MPTCP – Basic Use Cases



For seamless roaming

For high availability



MPTCP – Basic Use Cases





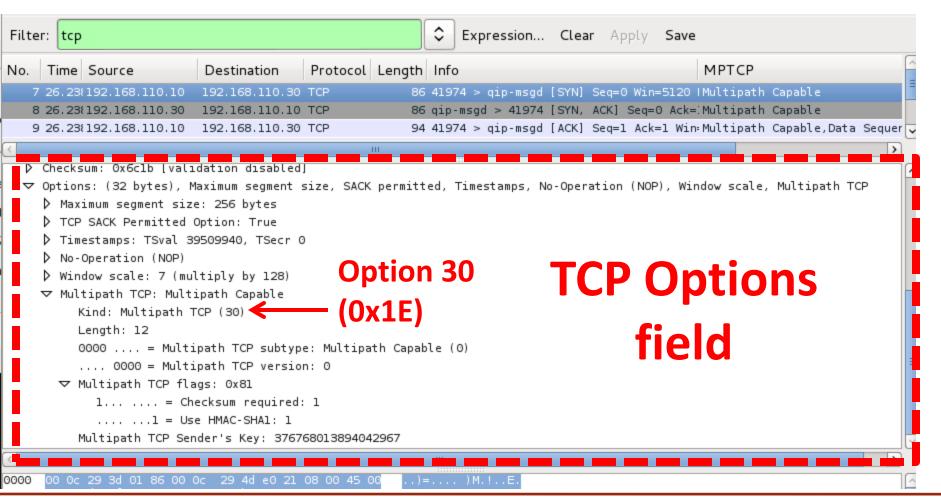
How is MPTCP implemented? – TCP Option

Bits 0 - 7		Bits 8 - 15	Bits 16 - 23		Bits 24 - 31	
	Sourc	e Port	Destination port			
		TCP Sequer	nce Number			
		TCP Acknowledgemer	nt Number (if	Ack Set)		
Data Offset	Reserved	TCP Flags (Ack, Syn etc)		Window Size		
	Chec	ksum		Urgent Point	er (if URG Set)	
Ox1e (MPTCP	Option Type)	Length	Subtype	MPTCP Ver	MPTCP Flags	
Remaining MPTCP Subtype Data						
Packet DATA						





Packet Breakdown - WireShark







How is MPTCP implemented? – MPTCP Subtypes

- 8 currently defined (ones relevant in bold)
- MP_CAPABLE Signals MPTCP support
- ADD_ADDR This address is also a way to reach me

 MP_JOIN - Add incoming subflow to the connection

- REMOVE_ADDR Please stop using [address] to reach me
- DSS How to map this stream's data against the overall data flow
- MP_PRIO
- MP_FAIL
- MP FASTCLOSE





Path Management - Linux

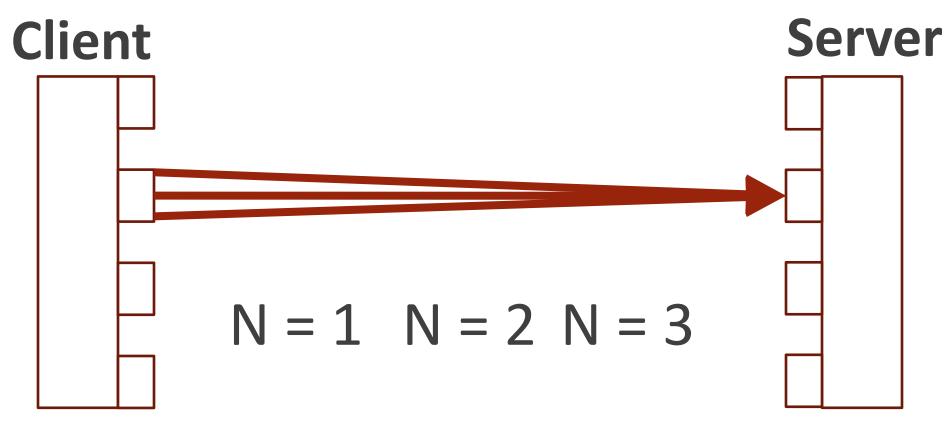
- The Linux Path Manager has two primary path managers at present
 - Fullmesh n:n (all to all)
 - Ndiffports 1-1 interfaces, n-1 ports

This is in the TCP stack... application layers get MPTCP for free (mostly)





Path Management - ndiffports



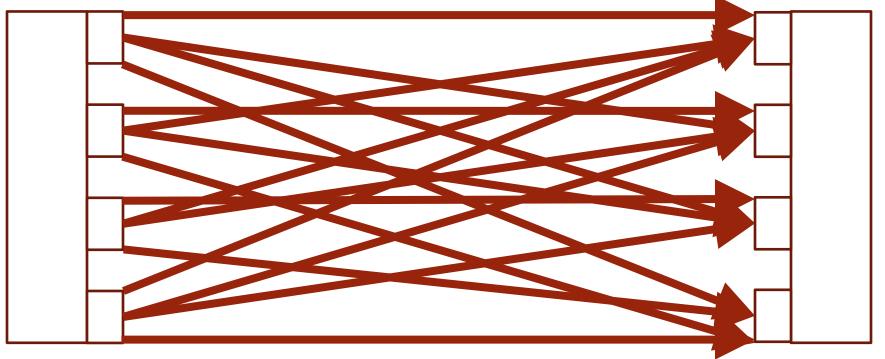
N different source ports,

1 destination port



Path Management - fullmesh

Client Server



All possible paths used





- TCP Handshake with additional details
- Data sequence numbering
 - Truncation of SHA1 of host key
- Authentication
 - MP_JOIN Challenge-response HMAC of other host's key, Nonce, AddressID
 - MP_FASTCLOSE Other party's key in plaintext
- Routing
 More on this later...
 - Packet sender decides which data goes down which path





MPTCP Cheatsheet

MPTCP Header:

Bits 0 - 7		Bits 8 - 15	Bits 16 - 23		Bits 24 - 31		
	Source Port			Destination port			
	TCP Sequence Number						
		TCP Acknowledgemer	nt Number (if	Ack Set)			
Data Offset	Reserved	TCP Flags (Ack, Syn etc)		Windo	w Size		
	Chec	ksum		Urgent Pointe	er (if URG Set)		
MP_C	apable	Length	Subtype	MPTCP Ver	MPTCP Flags		
Remaining MPTCP Subtype Data							
Packet DATA							

MPTCP	HEX	Flags?	Other Likely fields of interest
Subtype			
MP_CAPABLE	0x0		
MP_JOINs	0x1		
DSS	0x2		
ADD_ADDR	0x3		
REMOVE_ADDR	0x4		
MP_PRIO	0x5		
MP_FAIL	0x6		
MP_FASTCLOSE	0x7		

Getting the MPTCP Sequence Numbers:

Key	64 Bit number supplied by host	
Initial DSN (ISDN):	SHA1(key)[-64:]	Binary mode hash, network byte order

Initial DSS		
Subflow DSS	mapping likely starts at ISDN[0:32] + TCP ISN + 1	TCP Seq is 32 bits, + 1 for the
		SYN
MP_JOIN		

MP_JOIN Authentication (RFC 6824 Fig 8)

A		В			
TCP_SYN, MP_JOIN (TokenB, NonceA) ->					
	<- TCP_SYN_AC MP_JOIN(HM NonceB), NonceB)	K, IAC(Key=KB+KA, Msg = NonceA +			
TCP_ACK, MP_JOIN(HMAC(Key=KA+KB,Msg=NonceB+NonceA) ->					
<- TCP_ACK Token = ConnectionID = SHA1(Key)[0:32] of Other Party's key. (Capture from either steps 2 or 3 in the first handshake)					

Detecting MPTCP things

		Usage	Inbound	Detect inbound connection	TCP(SYN)
			Connection	attempts - Look for the SYN	TCP Option= 30 ** 00
			Attempts	packets with MPTCP Header	
			Successful	(Pre-viability) Look for Ack Packets	TCP(ACK)
			Handshake	with MPTCP Option header	TCP Option = 30 ** 00
			Valid	MPTCP Option header Look for	TCP(ACK)
			Handshake	Ack Packets with the MPTCP	TCP Option = 30 ** 00
Passive:				Option Header	
			MPTCP Joins	TCP SYN Packets with MPTCP TCP	TCP(SYN)
				Option and an MP_JOIN subtype	TCP Option = 30 ** 01
		Attacks	MPTCP Simple	Non look for non sequential last	



So who's using it?

- Nearly no one is using it large scale (yet), with a few exceptions
 - Apple iOS7 (Siri) → enabled by default in Yosemite (?)
 - Some other experimental stuff?
- Given that, there's a surprising number of implementations
 - Implementations available for several OS's (including Linux, BSD, Android), and baked in some way into commercial kit (Citrix, Cisco, Apple, Oracle, F5)
 - NOT Windows





Availability – Getting it working

- Linux
 - Linux reference implementation via apt-get (multipath-tcp.org) -- best way right now
 - Can work in Kali, but ... challenges

- Nicolas Maître made a ridiculously useful, near complete, SCAPY implementation
 - We're based some tools on this code, and fixed some bugs along the way
 - https://github.com/nimai/mptcp-scapy





Background **Technical Introduction Key Security Effects Perimeter Security Network Management MPTCP** Future





MPTCP's Key Security Effects

- Cross-path traffic fragmentation
 - That's the whole point!
- Moving target
 - Ability to change source and destination addresses in the middle of a connection
- Connection Resilience
 - Has additional checksums that require capture of the initial packet to reliably fake
 - Until every subflow is dead the overall connection keeps going
- "Reverse" connections



Because of these...

- Cross-path
- Moving target
- ConnectionResilience
- Reverse connections

... if your approach to security requires *any* of these...

- See all app layer data in a TCP stream
- Associate logical sessions to IP addresses
- Tamper with or close "bad" connections mid-stream
- Differentiate clients from servers based on connection direction

...then something is probably going to break





How practical are these attacks?

- Today? Extremely.
 - But only if both endpoints speak MPTCP
 - Of which... there aren't many. Yet.

- In an MPTCP world, a bit less
 - But we have to change the way we do things in network security





Practicality Going Forward

All of those things can be partially mitigated with MPTCP aware infrastructure and security tools.

But overall, there remain some interesting shifts in how network flows work – especially if we go in with "well meaning" intent



MPTCP's Key Security Effects

A few slides back...

The packet sender decides which data goes down which path.

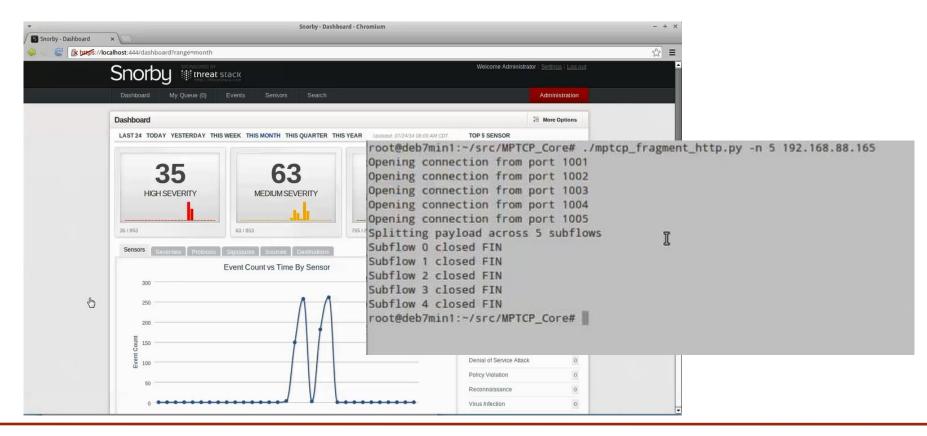
- Normal/benign clients won't choose pathological fragmentation schemes
 - But there's nothing stopping us...





PoC tool for MPTCP IDS Evasion

Demo!





Background **Technical Introduction Key Security Effects Perimeter Security Network Management MPTCP** Future

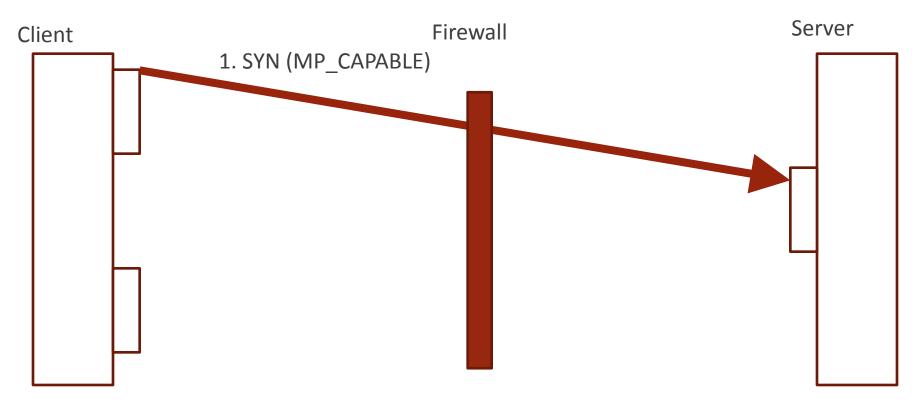




MPTCP changes things for perimeters

•How'd you like an outbound incoming connection?

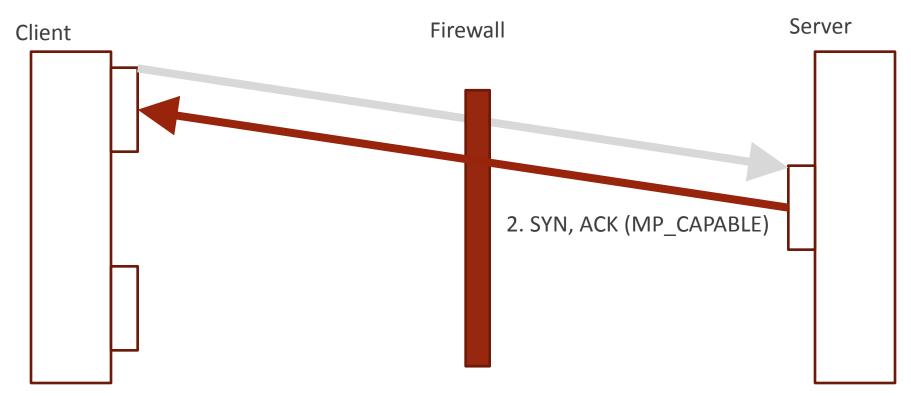




MPTCP connection looks like TCP so far...



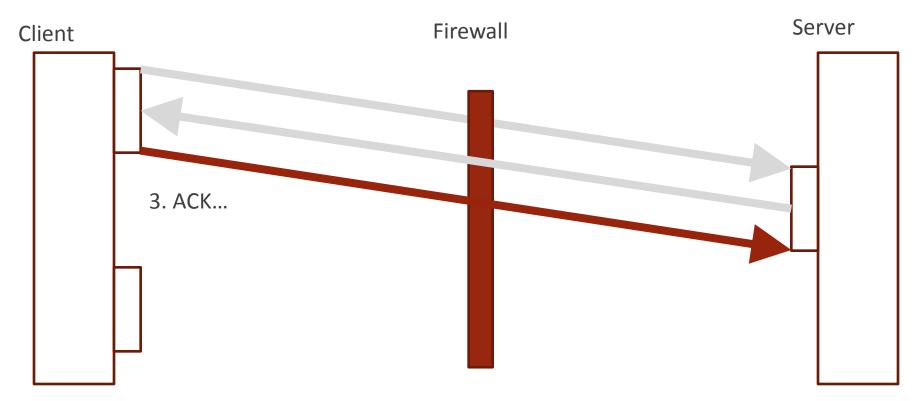




Still seems pretty standard, albeit with extra TCP OPTIONS



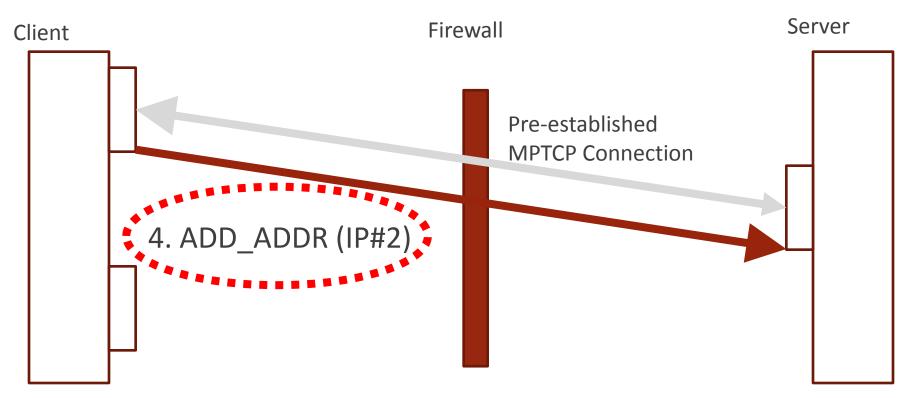




OK, so it's a TCP connection with an additional options... so what?

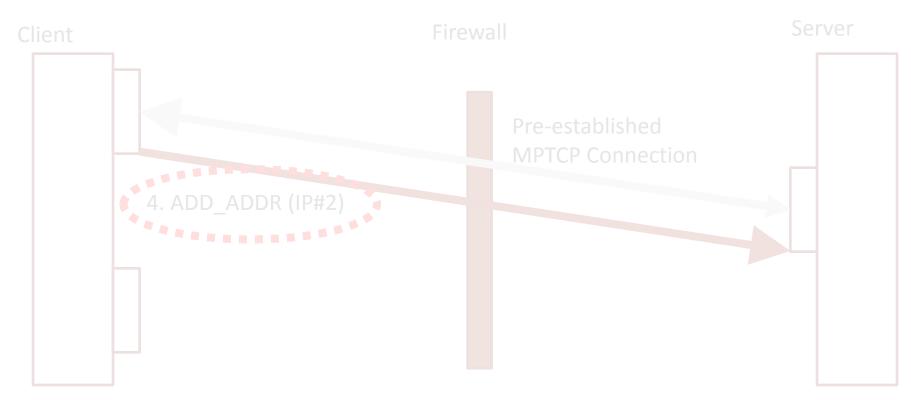






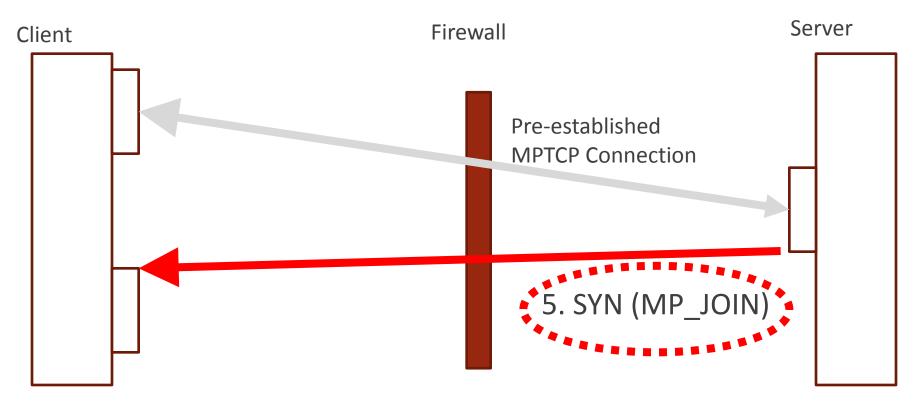
• Well, what if the client tells the server about a new address?





Now, the "Internal" host may set up a connection to the advertised address





Is this new connection incoming or outgoing?





Is this new connection incoming or outgoing?



Background **Technical Introduction Key Security Effects Perimeter Security Network Management MPTCP** Future





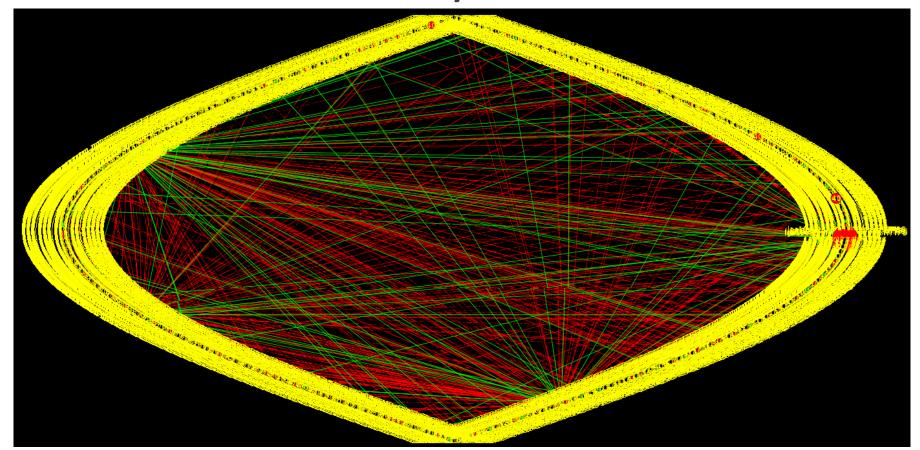
MPTCP and ... Network monitoring

If tool doesn't understand MPTCP, flows look like unrelated TCP streams





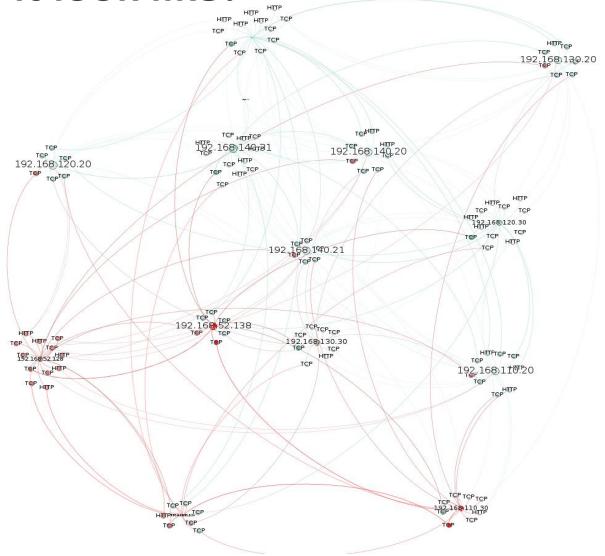
On the network: If you don't understand



Each yellow blob is actually part of an address label

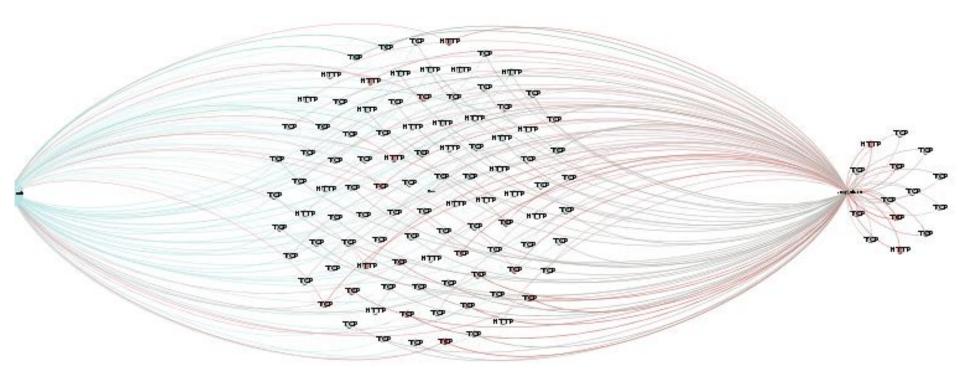


On the network: If you don't understand it, but you cluster IPs





On the network: If you do understand



But you can only do this when you can see & correlate all related flows...





MPTCP Defense - Awareness

- People
- Technology
 - Check support
 - Look for use
- Architecture
 - Terminate it where you terminate SSL



MPTCP Support Scanner

```
root@psthomas-neo-dev:~/mptools# ./scanner.py
usage: scanner.py [-h] [--ip SRC_IP] host port
Network scanner to test hosts for multipath TCP support. Requires root
privileges for scapy.
positional arguments:
               comma-separated IPs or ranges (globs allowed), eg
 host
               "127.0.0.1,192.168.1-254,203.0.113.*"
               comma-separated port(s) or port ranges, eg "22,80,8000-8999"
  port
```

```
root@psthomas-neo-dev:~/mptools# ./scanner.py 192.168.88.164 22,80
Testing: 192.168.88.164 ... on local network... at ARP: 00:0c:29:c8:8a:61
got MPTCP Response from 192.168.88.164 : 22 !... 20
RST Test indicates MPTCP support
qot MPTCP Response from 192.168.88.164 : 80 !... 20
RST Test indicates MPTCP support
****Results:****
       192.168.88.164
                        {22: 'MPTCP (MP_JOIN Verified)'}
                        {80: 'MPTCP (MP_JOIN Verified)'}
```





Accomplishes three things

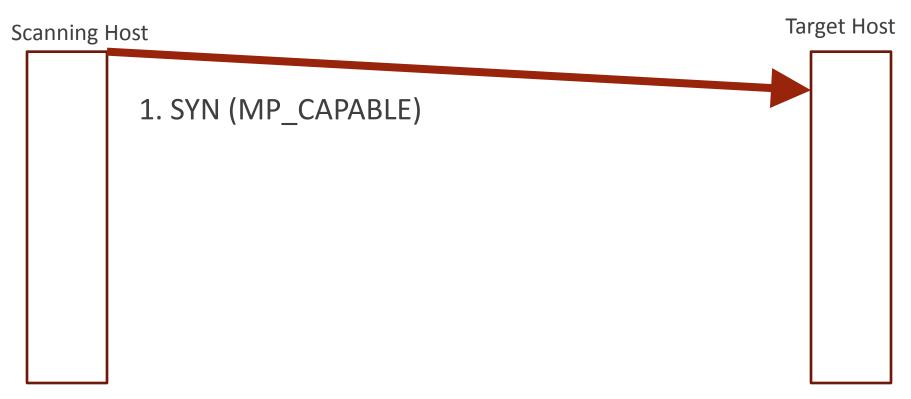
Test device for apparent support

 Test for actual support (as opposed to repeating the option blindly)

Test network path allows it to get there



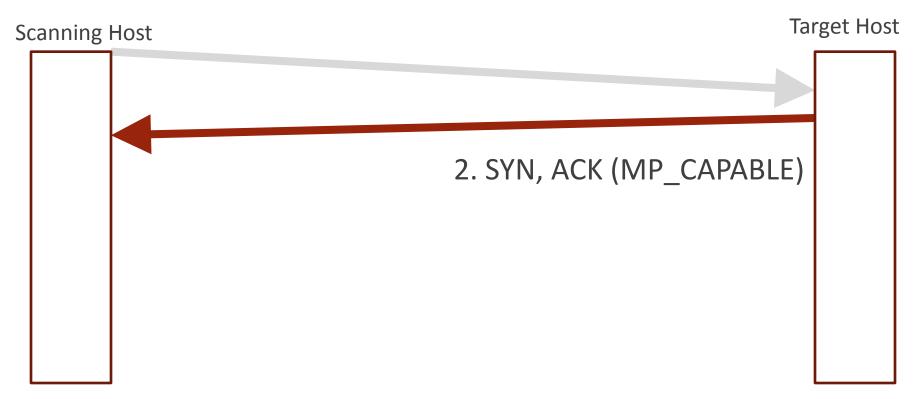




Send an MP_CAPABLE syn



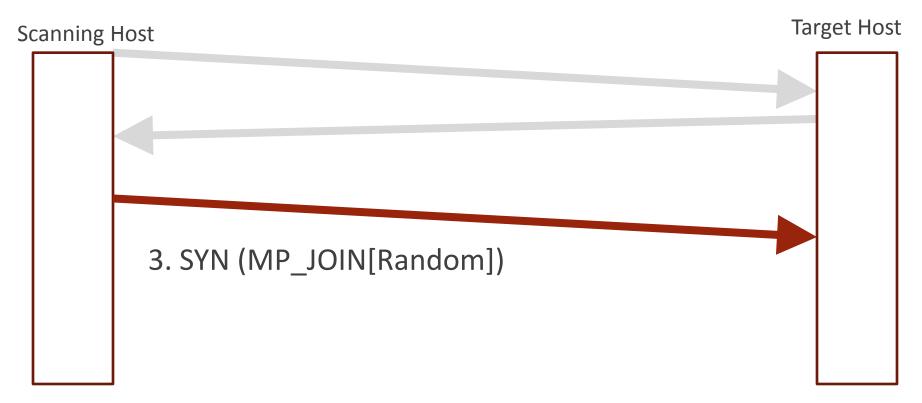




• We got an MP_CAPABLE response.. But is it genuine?







Send a join to an invalid connection ID







- An MPTCP host will RST an invalid join,
- An ACK reply indicates TCP only



MPTCP Stripping

- Transparent proxy on primary path
 - Either no MPTCP support, or only on the one interface

Firewall rules:

strip-options 30 - iptables,

tcp-options 30 30 clear - Cisco IOS





MPTCP and Active Network Security

To track & modify MPTCP, you must

1. Capture the initial handshake

- 2. Perform non-trivial calculations to determine
 - Connection membership
 - Correct checksum or modified traffic





Background **Technical Introduction Key Security Effects Perimeter Security Network Management MPTCP Future**



MPTCP and ... Privacy

 MPTCP shifts power towards endpoints, and away from infrastructure & ISP's

I don't trust my ISP or Cellular company...

But they probably don't trust each other either!





- What will change in a multipath future that simply cannot work with most existing security models?
 - Split trust crypto
 - Multipath agility

Some research into privacy effects already underway





 Multipath communications are awesome, and they're coming

 Multipath communication confounds business & security models relying on inspection

Now is the time for network security to prepare





PCAP Challenge (with solution) @ http://bit.ly/MPTCPTROOPERS15





Questions?

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Patrick Thomas @coffeetocode

pthomas@neohapsis.com

Downloads:

https://github.com/Neohapsis/mptcp-abuse

More stuff will be released @ http://labs.neohapsis.com

